# Using Recursively Defined Subsets of the Power Set of P to Show an Inequality betweenP andBPP 

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#### Abstract

In this paper we resolve the question of whether or not $P$ and BPP are unequal.We do this by showing the existence of a set that is a subset of $P$ and not in $P$, then using a property of the definition of this set to show it is within BPP.


## I. Introduction

In this paper we use the definition of BPP in terms of NP from 28 and use aproperty of this definition and proof of the existence of a certain class of setsand their non inclusion in P , followed by their inclusion in BPP, and the trivialproof that $P$ is an improper subset of BPP to show $P$ is a proper subset of BPP.

## II. Informal outline

We assume the definition of addition, subtraction and the complexity classes, however we assume a more abstract set of numerical definition as this allows forgreater expressive capability (i.e changing base without changing context). Weassume the definition of BPP in terms of NP seer28s. Our first theorem is that Pis an improper subset of BPP, this proof is trivial as if a is an element of $P$ thenthere exists a polynomial time deterministic turing machine that solves a. Whereobviously if there exists a polynomial time deterministic turing machine thatsolves a then there exists a polynomial time non deterministic turing machinesuch that if b is in a then at least two thirds of the computation paths areaccepting and if $b$ is not in a then less that one third of the computation pathsaccept. Our next theorem is that there exists a subset of $P$ defined using the setof elements in a (where $a$ is solvable by the polynomial time deterministic turingmachine $b$ ) such that every polynomial non deterministic accepting computationpaths in $c$ is in $d$ where a is a subset of $d$ and $d$ is in not within $P$ and the stateset input alphabet et al of $c$ is the same as that for $b$. The proof of this is simpleby constructing the functions $f 1$ and $f 2$ see section 6. Theorem 3 is the existenceof a subset of $P$ that is within BPP but not P. We prove this from the fact thatif the set we proved with theorem 2 is not exists then it is not within $P$ thenit is within BPP and we have obviously previously shown its existence. Our final theorem and its proof are trivial, that is P is a proper subset of BPP, wehave from theorem one that P is an improper subset of BPP and we have fromtheorem three that there exists a set in BPP that is not in P.

## III. Informal definitions and axioms

Definition. 1.

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\((a \oplus b)\) iff \((T(a) \neq T(b))\) where obviously \(T\) is the binary truth function of the
language
    Definition.2.
\(a \wedge b \operatorname{iff}((T(a)=T(b))(T(a)=1))\) where \(T\) is again the binary truth function
of the language
                                    Definition.3.
\(a \vee b \operatorname{iff}((a\) and \(b) \oplus(a \oplus b))\) Where in this section and \(b\) is equivalent to \(a \wedge b\)
                                    Definition. 4 .
\(a=\{b \mid c[d \mid e]\}\) iff f is in a iff d replaced by f in the statement c is true
                                    Definition.5.
set[a] iff there exists \(c\) and \(d\) such that \((a=\{b \mid c[d \mid b]\}))\) )
                                    Definition.6.
string[a,b] iff \(b\) is a set such that for all \(c\) and for all \(d\) such that \(c\) and \(d\) are in
\(\mathbf{b}(c d \oplus d c)\)
                                    Definition.7.
\(\operatorname{ath}(b, c)=d\) iff \(a<|c|\) where \(a\) is in IN and \(a=|\{e \mid(e \in c(e b d))\}|\)
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Definition. 8 .
$a=b$ iff for all $c, d$ such that string $[c, d]$ and $a$ is in $d$ where $c, d$ is true $c, d[a-b]$ is true for all $e, f$ such that string $[e, f]$ such that $b$ is in $f$ and $e, f$ is true $e, f[b-a]$ is true

Definition.9.
$a[b \mid c]=d$ iff a is a string and d is a string and there exists e such that the eth element of $a$ is equal to $b$ and the eth of $d$ is equal to $c$ such that for all $f$ where $f \neq e$ the fth element of a is equal to the fth element of d

Definition. 10.
$a / b=c$ iff for all $d$ where $d$ is in a and $d$ is not in $b$ we have $d$ in $c$
Definition 11.
class[a] iff there exists $b$ such that $b$ is in a such that $b$ is a set
Definition. 12.
$a \subseteq b$ iff if c is in a then c is in b and b is a set or class(where in this section a or b is equivalent to $a \vee b$ and a or mutually exlusive or os equivalent to $a \oplus b$

Definition. 13 .
$a \subset b$ iff if a is not equal to b and c is in a then c is in b

## Definition. 14.

$a \times b=c$ iff d is in c iff $\mathrm{d}=\mathrm{e}, \mathrm{f}$ and e is in a such that f is in b

## Definition. 15.

$S_{b}^{a}(c)=d$ iff $\mathrm{a}, \mathrm{b}$ is a string where c is in b and d is in b such that cad and there doesn't exists $f$ such that fad and caf

Definition. 16.
$a \simeq^{b, c} \mathrm{~d}$ iff for all $\mathrm{c}, \mathrm{d}$ such that $\mathrm{c}, \mathrm{d}$ is a string and a is in d where $\mathrm{c}, \mathrm{d}$ is true it is not true that if $c, d$ is true then $g b h$ where $g$ is in $c$ and $h$ is in $c$ however it is true that where $c, d[a-b]$ is true such that for all $e, f$ where $e, f$ is a string and $b$ is in $f$ and $e, f$ is true it is not true that if $e, f$ is true the where $i b j$ and $i$ is in $f$ and $\mathbf{j}$ is in $\mathbf{f}$ it is true that $e, f[b \mid a]$

Definition. 17.
$a=\{b, c\}(\mid \notin c(\operatorname{string}[b, c])))$ iff $d$ is in a iff $(d=\inf (b, c)$ or $d=\sup (b, c))$ or there exists e such that $e \simeq^{b, c}=$, and there exists f such that $f \simeq^{b, c}=$, where $\left(S_{c}^{b}(e)=d\right.$ and $S_{c}^{b}(d)=f$

Definition. 18.
$b, c$ is a string where $a$ is in $c$ and $a=b, c$ iff $(c=\{a\})$
Definition. 19.
$a c b=c, d$ where $c, d$ is a string iff $b$ is in $d$ and $a=c, e$ where $c, e$ is a string and $e=d /\{b\}$ such that for all $f$ such that $f$ is in $e$ we have $f c b$

Definition. 20.
$d e c i=\{0,1,2,3,4,5,6,7,8,9\}$
Definition. 21.
$<$, deci $=0<1<2<3<4<5<6<7<8<9$

Definition. 22.
$\mathbb{N}_{a, b}=\left\{a, c \mid\left(\operatorname{string}[a, c]\left(\forall d\left(d \in c\left(d \simeq^{a, c} e(e \in b)\right)\right)\right)\right\}\right.$ iff a,b is a string
Definition. 23.
$a: b \rightarrow c$ iff for all $d$ where $d$ is in $b$ there exists e such that e is in c and $\mathrm{a}(\mathrm{d})$ is equal to $e$ and for all $f$ such that $f$ is not in $b$ it is true that for all $g$ we have $a(f)$ is not equal to $g$

Definition. 24.
injection $[a, b, c]$ iff $a: b \rightarrow c$ and if $\mathrm{a}(\mathrm{d})=\mathrm{a}(\mathrm{e})$ then d is equal to e
Definition. 25.
surjection $[a, b, c]$ iff $a: b \rightarrow c$ where for all d such that d is in c there exists e where $e$ is in $b$ such that $a(e)=d$ )

Definition. 26.
bijection $[a, b, c]$ iff injection $[a, b, c]$ and surjection $[a, b, c]$
Definition.27.
$|a|<|b|$ iff there exists c such that injection $[c, a, b]$
Definition. 28 .
$|a|=|b|$ iff there exists c such that bijection $[c, a, b]$
Definition. 29 .
$a, b<\mathbb{N}_{a, c} a, d$ iff a,b is in $\mathbb{N}_{a, c}$ and a,d is in $\mathbb{N}_{c, d}$ and a,b is a string such that a, d is a string where $((|\boldsymbol{b}|<|d|) \oplus|b|=|d|$ such that there exists e such that for all f where $f<e$ and $\mathrm{fth}(\mathrm{a}, \mathrm{b})=\mathrm{fth}(\mathrm{a}, \mathrm{d})$ we have $\operatorname{eth}(\mathrm{a}, \mathrm{b})<_{\mathbb{N}_{a, c}}$ eth(a,d)

Definition. 30
$\mathbf{N}_{a, b}(c)=d$ iff $c$ is in $\mathbb{N}_{a, b}$ and $\left|\left\{e \mid\left(e<\mathbb{N}_{a, b} c\right)\right\}\right|=d$
Definition. 31 .
$\mathbb{N}=\left\{a \mid\left(\exists b\left(\exists c\left(\operatorname{string}[b]\left(\mathbb{N}_{b}(c)=a\right)\right)\right)\right)\right\}$
Definition. 32 .
$T M[a, b, c, d, e, f, g]$ iff c is in b and $d=b / c$ such that $g \subseteq a$ where f is in a and $e:(a / g) \times b \times\{T M L, T M R\} \rightarrow a \times b \times\{T M L, T M R\}$ and $|g|=\mathbb{N}_{<, \text {deci }}(2)$

Definition. 33 .
$a \lessgtr_{b, c} d$ iff a is a set or a class and $d$ is a set or a class and for all e such that e is in a and $e \simeq^{b, c} f$ it is true that f is in d

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\text { Definition. } 34 \text {. }
$$

compfunc $[a, b]$ iff , a is a string and for all c such that $\mathrm{b}(\mathrm{c})=\mathrm{d}(\mathrm{e}(\mathrm{c}))$ and compfunc $[\mathrm{f}, \mathrm{e}]$ where $\mathrm{d}=\sup (\mathrm{a})$ we have $f=a /\{d\}$

## Definition. 35 .

$\operatorname{DTM}[a, b, c, d, e, f, g, h, i, j]$ iff $, \mathrm{j}=\mathrm{A}, \mathrm{B}, \mathrm{C}$ and A is in g where B is in b and C is in IN such that $T M[a, b, c, d, e, f, g]$ and ,i is a string and, j is a string where $i \lesssim b$ and compfunc $[h, k]$ where $k\left(f,, i, \mathbb{N}_{<, \text {deci }}(0)\right)=, j$ and $|h|$ is in IN and $\left(\operatorname{lth}(, \mathrm{h})(\mathrm{m}, \mathrm{n}, \mathrm{o})=\mathrm{p}, \mathrm{q}, \mathrm{r}\right.$ iff $\mathrm{o}=\mathbf{N}_{<, \text {deci }}(0)$ and $\mathrm{e}(\mathrm{m}, \inf (\mathrm{n}))=\mathrm{p}, \mathrm{N}_{<, \text {deci }}(1)$, TML where $\mathrm{r}=\mathrm{N}_{<, \text {deci }}(0)$ and $\inf (\mathrm{q})=\mathrm{c}$ such that for all s such that $r<\mathbb{N} s$ where $\operatorname{sth}(n)=S_{\mathbb{I N}}^{\mathrm{⿺}}(s) \operatorname{th}(q) \oplus \operatorname{oth}(\mathrm{n})=\sup (\mathrm{n})$ and $\mathrm{e}(\mathrm{m}, \sup (\mathrm{n}))=\mathrm{p}, \operatorname{sth}(\mathrm{q})$, TMR where $S_{\mathrm{IV}}^{<}(0)=r$ and $\sup (\mathrm{q})=\mathrm{rth}(\mathrm{q})$ and $\mathrm{rth}(\mathrm{q})=\mathrm{c}$ such that for all t such that $\mathrm{t}<_{\mathrm{N}}$ o and $\operatorname{tth}(\mathrm{n})=\operatorname{tth}(\mathrm{q}) \oplus o \neq \mathrm{N}_{<, \operatorname{deci}}(0)$ and $\mathrm{e}(\mathrm{m}, \operatorname{oth}(\mathrm{n}))=\mathrm{p}, \mathrm{oth}(\mathrm{q})$, TML such that for all s such that $s \neq o$ it is true that $\operatorname{sth}(\mathrm{n})=\operatorname{sth}(\mathrm{q}) \oplus \operatorname{oth}(n) \neq \sup (n)$ and $(\mathrm{e}(\mathrm{m}, \operatorname{oth}(\mathrm{n}))=\mathrm{p}, \operatorname{oth}(\mathrm{q}), \mathrm{TMR}$ such that for all s such that $s \neq o$ it is true that $\operatorname{sth}(\mathrm{n})=\operatorname{sth}(\mathrm{q})$

Definition. 36 .
$N D T M[a, b, c, d, e, f, g, h, i, j]$ iff for all k such that k is in e it is true that $T M[a, b, c, d, k, f, g]$ and compfunc $[k, l]$ and $l\left(f, l, \mathbb{N N}_{<, \operatorname{dec} i}(0)\right)=, j$ where $, \mathbf{j}=\mathrm{A}, \mathrm{B}, \mathrm{C}$ such that $A$ is in $g$ and $B$ is in $b$ where $C$ is in $\mathbb{N}$ such that for all $m$ such that $\mathrm{m}<_{\mathbb{N}}|h|$ there exists n such that n is in e and $\mathrm{mth}(\mathrm{h})(\mathrm{o}, \mathrm{p}, \mathrm{q})=\mathrm{r}, \mathrm{s}, \mathrm{t}$ such that $\left(\left(\left(\mathrm{n}(\mathrm{o}, \mathrm{qth}(\mathrm{p}))=\mathrm{r}, \mathrm{qth}(\mathrm{s}), \mathrm{u}\right.\right.\right.$ and $q \neq \mathrm{N}_{<, \text {deci }}(0)$ such that $\left(u=T M L\left(S_{\mathrm{N}}^{\mathrm{N}}(t)=\right.\right.$ $q)) \oplus\left(q \operatorname{th}(p) \neq \sup (p)\right.$ andu $\left.\left.=T M R\left(S_{\mathbb{I N}}(q)=t\right)\right)\right) \oplus(n(o, q \operatorname{th}(p))=r, u t h(s), v$ and $u=\mathbb{N}_{<, d e c i}(1)$ such that $\inf (\mathrm{s})=\mathrm{c}$ such that for all w such that $w t h(p)=$ $S_{\mathrm{IN}}^{\mathbb{<}}(w) \operatorname{th}(s)$ and $\mathrm{v}=\mathrm{TML}$ where $\left.\left(S_{\mathrm{IN}}^{<}(t)=q\right)\right) \oplus \mathrm{u}=\mathrm{q}$ and $S_{\mathrm{N}}^{<}(q) \operatorname{th}(s)=\sup (s)$ such that $q \operatorname{th}(p)=\sup (p)$ and $v=T M R$ and $\sup (s)=c$ such that for all $w$ such that $w<_{\mathbb{N}} q$ and $w \operatorname{th}(\mathrm{p})=\mathrm{wth}(\mathrm{s})$

## Definition. 37.

$a=b^{c}$ where $a \in \mathbb{I N}$ iff compfunc[d,e] and $e\left(\mathbb{N}_{<, \text {deci }}(1)\right)=a$ such that $|d|=c$ such that for all f such that $\operatorname{fth}(, d)(g)=$ in $(b, g)$

## Definition. 38 .

$a(b)=O_{c}(d(b))$ iff there exists e and there exists f such that e is in $\mathbb{I N}$ and f is in IN such that for all g such that g is in c and $a(g) \leqslant \mathrm{iv} \mathbb{N}\left(e, d(g)^{f}\right)$

$$
\text { Definition. } 39
$$

$|a| \leqslant|b|$ iff $((|a|=|b|) \oplus(|a|<|b|))$
Definition. 40 .
$+_{\mathbb{N}}(a, b)=c$ iff $c$ is in $\mathbb{N}$ and a is in $\mathbb{N}$ and b is in $\mathbb{N}$ and $\left\{d \mid\left(d<_{\mathbb{N}} c\left(a \leqslant_{\mathbb{N}} d\right)\right)\right\}$
Definition. 41 .
in $(a, b)=c$ iff c is in $\mathbb{N}$ and a is in $\mathbb{N}$ and b is in $\mathbb{N}$ such that compfunc[d, e$]$ where $|d|=b$ and $e\left(\mathbf{N}_{<, \text {deci }}(0)\right)=c$ such that for all f such that for all g it is true that $\operatorname{fth}(d)(g)=+\mathbb{N}_{<\text {dect }}(a, g)$

Definition. 42 .
$P=\{k \mid((\operatorname{set}[k] \vee$ class $[k])(\exists a(\exists b(\exists c(\exists d)(\exists e(\exists f(\exists g(\exists l) \exists m)(\forall i(\exists h(\exists j) N D T M[a, b, c, d, e, f, g, h, i, j](|h|=$ iv $\left(l,|i|^{m}\right)(j=n, o, p(n=\inf (, g) \Leftrightarrow i \notin k)((n=\sup (, g) \Leftrightarrow i \in k)(n=$


Definition. 43 .
$N P=\left\{k \mid(\exists a)\left(\exists b\left(\exists c(\exists d)\left(\exists e\left(\exists f(\exists g)(\exists l(\exists m) \quad \forall i), i \in k\left(\exists h(\exists j) N D T M[a, b, c, d, e, f, g, h, i, j]\left(|h|={ }_{\mathrm{i}}\right.\right.\right.\right.\right.\right.\right.$ $\left(l,|i|^{m}\right)(j=n, o, p(n=\inf (, g) \Leftrightarrow i \notin k)((n=\sup (, g) \Leftrightarrow i \in k)(n=$


Defintion. 45.
$\mathbb{Z}_{a, b}=\left\{a, c \mid\left(\operatorname{string}[a, c]\left(a, c \in \mathbb{N}_{a, b} \oplus\left(a, c=-\operatorname{ad}\left(\operatorname{string}[a, d]\left(a, d \in \mathbb{N}_{a, c}\right)\right)\right)\right)\right)\right\}$
Defintion. 46.

$$
\mathbb{Z}_{a}(b)=c \Leftrightarrow\left(\mathbf{N}_{a}(b)=c \oplus\left(b=-\operatorname{de}\left(\operatorname{string}[d, e]\left(c=-\mathbf{N}_{a}(d, e)\right)\right)\right)\right)
$$

Defintion. 47.

$$
\mathbb{Z}=\left\{a \mid\left(\exists b\left(\exists c\left(\mathbb{Z}_{b}(c)=a\right)\right)\right)\right\}
$$

Defintion. 48.
$\mathbb{R}_{a, b}=\left\{a, c \mid\left(\right.\right.$ string $[a, c]\left(\left(a, c \in\right.\right.$ mathds $Z_{a, b} \oplus\left(. \in c\left(\forall e\left(e \in c\left(\sup (a, c) \neq\left(\forall g\left(g \in c\left(g \simeq^{a, c} h(g \neq .(h \in b\right.\right.\right.\right.\right.\right.\right.$ $\left.\left.\left.)))))))))) \oplus\left(\inf (a, c)=-\left(\sup (a, c) \neq .\left(\forall g\left(g \in c\left(g \neq-\left(g \neq .\left(g \simeq^{a, c} h(h \in h(h \in b))\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\} \Leftrightarrow \operatorname{string}[a, b]$

Defintion. 49 .

$$
\mathbb{R}_{a}(b)=\mathbb{Z}_{a}(b) \Leftrightarrow b \in \mathbb{Z}_{a}
$$

Definition.50.
${\stackrel{+}{\mathbb{Z}_{a, b}}}(a, c, a, d)=a, e \Leftrightarrow\left(\left({\stackrel{+}{\mathbf{Z}_{a, b}}}(a, c, a, d)=a, e \oplus\left({\stackrel{+}{\mathbb{Z}_{a, b}}}(a, c /\{\inf (a, c)\}, a, d /\{\inf (a, d)\})=a, e\right.\right.\right.$
$\left.(\inf (a, c)=-(\inf (a, d)=-)))) \oplus\left(\inf (a, d)=-\left(\inf (a, c) \neq-\left(-\mathbb{N}_{a, b}(a, c, a, d /\{\inf (a, d)\})\right)\right)\right)\right)$
Defintion. 51.
$\mathbb{R}=\left\{a \mid\left(\exists d\left(\exists b, c\left(\operatorname{string}[b, c]\left(\mathbb{R}_{b, c}(d)=a\right)\right)\right)\right)\right\}$

Definition. 52.

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\mp@subsup{\stackrel{~}{\mathbb{R}}}{a,b}{\prime}
((\neg(\existsf(fth(c)=.)))(gth(d)=.(c=a,i(d=a,j(\mp@subsup{+}{\mp@subsup{\mathbb{Z}}{a,b}{\prime}}{\prime}(c,{l|(l=mth(d)(m<N g))})=n(
e=a,n\cup{m|(m=oth(d)(g\mp@subsup{\leqslant}{\mathbb{N}}{}o))}(sup(a,n)a(a,{m|(m=oth(d)(g\mp@subsup{\leqslant}{\mathbb{N}}{}o))}))))))))))
\oplus(((inf(c) =- (inf(d)=-))\oplus(inf(c)\not=-(inf(d)\not=-)))(fth(c)=.(gth(d)=.(f\leqslant\mathbb{N}g(\operatorname{sup}(c)=jth(c)
(sup (d) = kth(d)(compfunc[l,m] (m}(d)=e(nth(,l)(d)=o\Leftrightarrow(((n<\mathbb{N}\mp@subsup{|}{\mathbb{N}}{}(k,f
(invnth (o) = + +
    (\forallp(p\not=n(invpth(o)=invpth(d))))))\oplus(n<\mp@subsup{<}{\mathbb{N}}{}\mp@subsup{-}{\mathbb{N}}{}(k,g)(-\mp@subsup{|}{\mathbb{N}}{}(j,f)<\mp@subsup{<}{\mathbb{N}}{}n
(invnth (o) = +}\mp@subsup{\stackrel{N}{N}}{a,b}{+}(\operatorname{invS}\mp@subsup{S}{\mathbb{N}}{<}nth(c),\operatorname{invnth}(d))(\mp@subsup{\stackrel{~}{\mathbb{N}}}{a,b}{}(\operatorname{invS}\mp@subsup{S}{\mathbb{N}}{<}nth(c),\operatorname{invnth}(d))=a,
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    \oplus(n<<\mathbb{N}-\mathbb{N}
(invS}\mp@subsup{S}{\mathbb{N}}{<}(n)\operatorname{th}(c),invnth(d)))(\mp@subsup{\stackrel{\mathbb{N}}{a,b}}{+}{~
    (\forallp(p\not=n(invpth(o)=\operatorname{invpth}(\mp@subsup{\stackrel{\mathbb{N}}{a,b}}{+}{(inf(a,q),d)))))))))))))))))))))})
                                    Defintion.53.
\stackrel{+}{\mathbb{R}}
                    Definition.54.
                    \mathbb{R}}
                    Definition.55.
    Z Z
\oplus(a,c\in\mp@subsup{\mathbb{N}}{a,b}{}(\operatorname{inf}(a,d)=-(\mp@subsup{\textrm{N}}{a,b}{}(a,c,a,d/{\operatorname{inf}(a,d)})=a,f(a,e=-aa,e)))))
    \oplus(inf(a,c)=-(inf(a,d)=-(\mp@subsup{\mathbb{N}}{a,b}{}(a,c/{\operatorname{inf}(a,c)},a,d/{\operatorname{inf}(a,d)})))))
                    Definition.56.
Z}(a,b)=c\leftrightarrow(\existsd(\existse(\existsf(\existsg(\mp@subsup{\mathbb{Z}}{d}{}(e,f)=g(\mp@subsup{\mathbb{Z}}{d}{}(e)=a(\mp@subsup{\mathbb{Z}}{d}{}(f)=b(\mp@subsup{\mathbb{Z}}{d}{}(g))))))))
                    Defintion.57.
                    / Z
                            Defintion.58.
                /ZZ}(\boldsymbol{a},\boldsymbol{b})=c\leftrightarrow(i\mathbb{Z}(\boldsymbol{c},\boldsymbol{b})=\boldsymbol{a}
                            Defintion.59.
        (\mp@subsup{\dot{\mathbf{R}}}{a,b}{}}(a,c,a,d)=a,e(a,c\in\mp@subsup{\mathbb{Z}}{a,b}{}\oplusa,d\in\mp@subsup{\mathbb{Z}}{a,b}{}))\Leftrightarrow((h(i)=j)\Leftrightarrow(i\in
            (iasup}(a,b)(\mp@subsup{S}{\mp@subsup{\mathbf{R}}{a,b}{<}}{<\mp@subsup{\mathbf{R}}{a,b}{}}(\operatorname{sup}(a,b))=k(\mp@subsup{N}{}{\mp@subsup{N}{a,b}{}}(i,j)=k))))
            (\dot{\mathbf{Z}}
        m(l))=a,e((qth(,m)(r)=s\Leftrightarrows=/\mp@subsup{\mathbf{Z}}{a,b}{}(r,h(\stackrel{\mathbf{N}}{+}(q,f)th(a,d)))))))))
                    Defintion.60.
    (\mp@subsup{\dot{R}}{a,b}{}}(a,c,a,d)=a,e(\neg(a,c\in\mp@subsup{\mathbb{Z}}{a,b}{}\oplusa,d\in\mp@subsup{\mathbb{Z}}{a,b}{}))))\Leftrightarrow(((\mp@subsup{\dot{\mathbf{Z}}}{a,b}{}(a,c,a,d)=a,e)
            (fth(a,c) =.(gth(a,d) =. (((h(i)=j)\Leftrightarrow(i\inb
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(\mp@subsup{\dot{Z}}{a,b}{}(a,{l|(l=mth(a,c)(m<f))},a,{l|(l=mth(a,d)(m<g))})=l(compfunc[m,n](compfunc[o,p]
(o(m(l)) =a,e((qth(,m) (r) =s\Leftrightarrows=/\mp@subsup{\mathbf{z}}{a,b}{}(r,h(\stackrel{+}{\mathbf{N}}=\mp@code{q}(q,f)\operatorname{th}(a,d)))(qth(,o)(r)=s
            \Leftrightarrows=/\mp@subsup{\mathbf{Z}}{a,b}{}(r,h(+
                    Definition.61.
    BPP}={a|(\existsx(\existsy(x\in\mathbb{R}(y\in\mathbb{R}(,c,d,e,f,g,h,(\foralli(j(i)=|{l|(\existsm
            NDTM[b,c,d,e,f,g,h,l,k,m](|l|\leqslant\mp@subsup{\leqslant}{\textrm{R}}{}\textrm{R}
(((k\ina) => (汭}(j(k),\mp@subsup{\mathbb{R}}{<,\mathrm{ deci }}{}(/\mp@subsup{\mathbf{R}}{,,\mathrm{ deci }}{}(2,3))\leqslant\mathbf{R}}|{l|(\existsm(NDTM[b,c,d,e,f,g,h,l,k,m
(m=n,o,p(n=\operatorname{sup}(,h)(|l|\leqslant\mathbb{R}\mathbb{R}(x,|k\mp@subsup{|}{}{y}))))))}|)))((k\not\ina(k\in\mp@subsup{c}{}{*}))=>(|{l|(\existsm(NDTM[b,c,d,e,f,g,h,l,k,m]
(m=n,o,p(n=\operatorname{sup}(,h)(|l|\mp@subsup{\leqslant}{\mathbb{R}}{}\mathbb{R}(x,|k\mp@subsup{|}{}{l}))))))}\mp@subsup{\leqslant}{\mathbb{R}}{}\mathbf{R}(\mp@subsup{\mathbb{R}}{<,\operatorname{deci}}{}(/\mp@subsup{\mathbb{R}}{<,\operatorname{deci}}{}(1,3)),j(k))))))))))))}
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                                    Defintion.62.
a,bDTM[c,d,e,f,g,h,i]}\Leftrightarrow(\forallj(\existsk(\existsl)DTM[c,d,e,f,g,h,i,k,j,l](|k|\leqslantb(|j|(l=m,n,
                        ((m=sup(,i)\Leftrightarrowj\ina)
        (m=inf(,i)\Leftrightarrow(j\not\ina))))))))))
                                    Definition.63.
(a}\mp@subsup{|}{}{b}(b\in\mathbb{NN}(a\in\mathbb{N})))\Leftrightarrow(\operatorname{compfunc}[c,d](|d|=b(\foralle(eth(,c)(f)=iv (f,a)(d(\mathbb{N}<,\operatorname{deci}(1))))))
                                    Defintion.64.
    a\in polyfunc \Leftrightarrow(\existsd(\existse(\forallb(a(b)=c\Leftrightarrow(c=\mp@subsup{|}{\mathrm{ iv }}{}(d,\mp@subsup{b}{}{e}))))))
                            Axiom.1.
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there exists a such that $a$ is a set
Axiom. 2.
there exists a such that $a$ is a class
Axiom. 3.

## $P \neq \varnothing$

## Axiom. 4

there exists a such that $a=\mathbb{N}_{<, \text {deci }}$
Axiom.5.
there exists a such that $a=\mathbb{N}$

## IV. Informal Theorems and proof

## $T h .1 . P \subseteq B P P$

$\operatorname{Pr.1.}((a \in P) \Rightarrow(\exists b, c, d, e, f, g, h(a, j D T M[b, c, d, e, f, g, h](j \in \operatorname{polyfunc}))))$
$\wedge(((\exists b, c, d, e, f, g, h(a, j D T M[b, c, d, e, f, g, h](j \in$ poly func $))) \Rightarrow(\exists x(\exists y(x \in \mathbb{R}(y \in \mathbb{R}(\exists b, c, d, e, f, g, h,(\forall i)$

$$
j(i)=\left|\left\{l \mid\left(\exists m\left(N D T M[b, c, d, e, f, g, h, l, k, m]\left(|l| \leqslant_{\mathbb{R} R}\left(x,|k|^{y}\right)\right)\right)\right)\right\}\right|
$$

$\left(\left((k \in a) \Rightarrow\left(\dot{\mathbf{R}}\left(j(k), \mathbb{R}_{<, \text {deci }}\left(/ \mathbf{R}_{,, \text {deci }}(2,3)\right) \leqslant_{\mathbf{R}} \mid\{l \mid(\exists m(N D T M[b, c, d, e, f, g, h, l, k, m]\right.\right.\right.\right.$
$\left.\left.\left.\left.\left.\left(m=n, o, p\left(n=\sup (, h)\left(|l| \leqslant_{\mathbb{R} R}\left(x,|k|^{y}\right)\right)\right)\right)\right)\right\} \mid\right)\right)\right)\left(\left(k \notin a\left(k \in c^{*}\right)\right) \Rightarrow(\mid\{l \mid(\exists m(N D T M[b, c, d, e, f, g, h, l, k, m]\right.$ $\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left(m=n, o, p\left(n=\sup (, h)\left(|l| \leqslant_{\mathbb{R}} \mathbf{R}\left(x,|k|^{l}\right)\right)\right)\right)\right)\right)\right\} \leqslant_{\mathbb{R}} \mathbb{R}\left(\mathbb{R}_{<, \operatorname{deci}}\left(/ \mathbb{R}_{<, \operatorname{deci}}(1,3)\right), j(k)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\}$
Th.2. $(\exists a(a \subseteq P($ string $[, a]($ string $[, b](|a|=|b|(\forall j(\forall k(\exists l(D T M[c, d, e, f, g, h, i, j t h(, b), k, l]$ $(l=m, n, o((m=\sup (, i) \Leftrightarrow(k \in j t h(, a)))((m=\inf (, i) \Leftrightarrow(k \notin j \operatorname{th}(, a)))(\forall j(N D T M[c, d, e, f, b, h, i, j, k, l]$ $\left.\left(|j| \leqslant_{\mathbf{R} \mathbf{R}}\left(\mathbb{R}_{<, d e c i}(10),|k|^{\mathbb{R}_{<, d e c i}(100)}\right)\right)\right)(\{k \mid(k \in a \vee(N D T M[c, d, e, f, b, h, i, j, k, l]$

$$
(l=m, n, o(m=\sup (, i)))))))\} \notin P())))))))))))
$$

Pr.2. $\left(\left((((a) b, c)=d \Leftrightarrow b t h(, c)=d)\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b, d D T M[e, f, g, h, y t h(, i), j, k]\right.\right.\right.$
$(d \in \operatorname{polyfunc}(\forall l(l \in a(\exists m(m \in \operatorname{polyfunc}(\exists n(n \in q(l, m D T M[e, f, g, h, n, j, k](\forall o(\forall p(p \in \operatorname{poly} f u n c($
$\left.\left.\left.o, p N D T M[e, f, g, h, q, j, k](\forall r(r \in o(r \in a))))))))))))))))))\}\left(\{01\} \in a_{1}\left(a_{1} \in P\right)\right)\right)\right) \Rightarrow(a \in a)\right) \wedge$
$\left(\left(f_{1}(d, 0)=0,0, \operatorname{TMR}\left(f_{1}(0,1)=\sup (, h), 1, \operatorname{TMR}\left(f_{2}(d, 1)=0,0,, \operatorname{TMR}\left(f_{2}(0,0)=\sup (, h), 0, T M R\right.\right.\right.\right.\right.$
$\left(f_{1}(d, 1)=2,0, \operatorname{TMR}\left(f_{1}(2,0)=\inf (, h), 0, \operatorname{TMR}\left(f_{1}(2,1)=\inf (, h), 0, \operatorname{TMR}\left(f_{2}(0,1)=\inf (, h), 0, T M R(\right.\right.\right.\right.$
$f_{2}(2,0)=\inf (, h), 0, \operatorname{TMR}\left(f_{2}(2,1)=,\inf (, h), 0, \operatorname{TMR}\left(f_{2}(d, 0)=\inf (, h), 0, \operatorname{TMR}(\right.\right.$
$f_{1}:\{0,1,2\} X\{0,1\} \rightarrow\{0,1,2\} X\{0,1\} X\{T M L, T M R\}\left(f_{2}:\{0,1,2\} X\{0,1\} \rightarrow\{0,1,2\} X\{0,1\} X\{T M L, T M R\}\right.$ $))()))))))))) \Rightarrow\left(\left((a(b, c)=d \Leftrightarrow b \operatorname{th}(, c)=d)\left(\exists a_{1}\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(\right.\right.\right.\right.$
$b, d D T M[\{0,1,2\},\{0,1\}, g, h, i, j, k](d \in \operatorname{polyfunc}(\forall l(l \in a(\exists m(m \in \operatorname{poly} f u n c($
$\exists n(n \in q(l, m D T M[\{0,1,2\},\{0,1\}, g, h, n, j, k](\forall o(\forall p(p \in \operatorname{polyfunc}(o, p N D T M[\{0,1,2\},\{0,1\}, g, h, q, j, k]$
$\left.\left.\left.\left.\left.(\forall r(r \in o(r \in a))))))))))))))))))\}\left(\{01\} \in a_{1}\right)\right)\right)\right)\right)\right) \wedge((a(b, c)=d \Leftrightarrow b \operatorname{th}(, c)=d)($
$\exists a_{1}\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b, d D T M[\{0,1,2\},\{0,1\}, g, h, i, j, k](d \in \operatorname{polyfunc}(\forall l(l \in a(\exists m)\right.$
$m \in \operatorname{polyfunc}(\exists n(n \in q(l, m D T M[\{0,1,2\},\{0,1\}, g, h, n, j, k]$
$\left.\left.(\forall o(\forall p(p \in \operatorname{polyfunc}(o, p N D T M[\{0,1,2\},\{0,1\}, g, h, q, j, k](\forall r(r \in o(r \in a)))))))))))))))))))\}\left(\{01\} \in a_{1}\right)\right)\right) \Rightarrow$ $\left(\left(\neg\left(\left((a(b, c)=d \Leftrightarrow b \operatorname{th}(, c)=d)\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b \subseteq y(b, d D T M[e, f, g, h, i, j, k]\right.\right.\right.\right.\right.$ (d polyfunc ( $\forall l\left(l \in a_{1}(\exists m(m \in \operatorname{polyfunc}(\exists n)(n \in q(l, m D T M[e, f, g, h, n, j, k](\forall o(\forall p(p \in \operatorname{polyfunc}(\right.$ $\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.o, p N D T M[e, f, g, h, q, j, k]\left(\forall r\left(r \in o\left(r \in a_{1}\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\}\left(\left(\{01\} \in a_{1}\left(a_{1} \in P\right)\right)\right)\right) \Rightarrow$ $\left((a(b, c)=d \Leftrightarrow b \operatorname{th}(, c)=d)\left(\forall a_{1}\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b, d D T M[e, f, g, h, i, j, k]\right.\right.\right.$ ( $d \in$ polyfunc $(\forall l(l \in a(\exists m(m \in \operatorname{polyfunc}(\exists n(n \in q(l, m D T M[e, f, g, h, n, j, k])(\forall o(\forall p(p \in$ polyfunc $($ $\left.\left.\left.\left.\left.\left.o, p N D T M[e, f, g, h, q, j, k](\forall r(r \in o(r \in a))))))))))))))))))\}\left(\{01\} \in a_{1}\left(a_{1} \notin P\right)\right)\right)\right)\right)\right)\right)\right)$ $\wedge\left(\left(\neg\left((a(b, c)=d \Leftrightarrow b \operatorname{th}(, c)=d)\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b \subseteq y(b, d D T M[e, f, g, h, i, j, k](d \in\right.\right.\right.\right.$ polyfunc $\left(\forall l\left(l \in a_{1}(\exists m(m \in \operatorname{polyfunc}(\exists n)(n \in q(l, m D T M[e, f, g, h, n, j, k])(\forall o(\forall p(p \in \operatorname{polyfunc}(\right.\right.$ $\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.o, p N D T M[e, f, g, h, q, j, k]\left(\forall r\left(r \in o\left(r \in a_{1}\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\}\left(\{01\} \in a_{1}\left(a_{1} \in P\right)\right)\right)\right)\right) \Rightarrow$ $\left((a(b, c)=d \Leftrightarrow b t h(, c)=d)\left(\forall a_{1}\left(a_{1}=\{a(z, b) \mid(z \in \mathbb{N}(b \in P(b, d D T M[e, f, g, h, i, j, k](d \in\right.\right.\right.$ polyfunc ( $\forall l(l \in a(\exists m(m \in \operatorname{polyfunc}(\exists n(n \in q(l, m D T M[e, f, g, h, n, j, k](\forall o(\forall p(p \in$ polyfunc $($ $\left.\left.\left.\left.o, p N D T M[e, f, g, h, q, j, k](\forall r(r \in o(r \in a))))))))))))))))))\}\left(\{01\} \in a_{1}\left(a_{1} \notin P\right)\right)\right)\right)\right)\right)$ Th. $3 .(\exists a(a \subseteq P(a \notin P(a \in B P P))))$
$\operatorname{Pr} .3(a \subseteq P($ string $[, a]($ string $[, b](|a|=|b|(\forall j(\forall k(\exists l(D T M[c, d, e, f, g, h, i, j t h(, b), k, l]$
$(l=m, n, o((m=\sup (, i) \Leftrightarrow(k \in j \operatorname{th}(, a)))((m=\inf (, i) \Leftrightarrow(k \notin j \operatorname{th}(, a)))(\forall j(N D T M[c, d, e, f, b, h, i, j, k, l]$ $\left.\left(|j| \leqslant \mathbb{R}_{\mathbf{R}}\left(\mathbb{R}_{<, \text {deci }}(10),|k|^{\mathbb{R}_{<, \text {deci }}(100)}\right)\right)\right)(\{k \mid(k \in a \vee(N D T M[c, d, e, f, b, h, i, j, k, l]$
$(l=m, n, o(m=\sup (, i)))))))\}=z(z \notin P)))))))))))))) \Rightarrow(z \in B P P)) \wedge(\exists a(a \subseteq P(\operatorname{string}[, a](\operatorname{string}[, b]$
$(|a|=|b|(\forall j(\forall k(\exists l(D T M[c, d, e, f, g, h, i, j \operatorname{th}(, b), k, l]$
$(l=m, n, o((m=\sup (, i) \Leftrightarrow(k \in j \operatorname{th}(, a)))((m=\inf (, i) \Leftrightarrow(k \notin j \operatorname{th}(, a)))(\forall j(N D T M[c, d, e, f, b, h, i, j, k, l]$
$\left.\left(|j| \leqslant_{\mathbf{R}} \mathbf{R}\left(\mathbb{R}_{<, \text {deci }}(10),|k|^{\mathbf{R}<, \text { deci }}(100)\right)\right)\right)(\{k \mid(k \in a \vee(N D T M[c, d, e, f, b, h, i, j, k, l]$
$(l=m, n, o(m=\sup (, i)))))))\}=z(z \notin P)))))))))))))))$
Theorem. 4 .
$P$ is a proper subset of BPP
Proof.4.
There exists a set in BPP that is not in P and we already have from theorem 1 that $P$ is a improper subset of BPP
V. Formal de_nitions, axioms and proof
$D .1 .(a \oplus b) \leftrightarrow(T(a) \neq T(b))$
$D .2 . a \wedge b \Leftrightarrow((T(a)=T(b))(T(a)=1))$
$D .3 . a \vee b \leftrightarrow((a \wedge b) \oplus(a \oplus b))$
D.4.a $=\{b \mid c[d \mid e]\} \Leftrightarrow(\forall f(c[d \mid f](c[d \mid f] \in a)))$
$D .5 \cdot \operatorname{set}[a] \leftrightarrow(\exists c(\exists d(a=\{b \mid c[d \mid b]\})))$
D.6.string $[a, b] \leftrightarrow(\operatorname{set}[b](\forall c(\forall d(c \in b(d \in b(c d \oplus d c))))))$
$D \cdot 7 \cdot \operatorname{ath}(b, c)=d \leftrightarrow(a<|c|(d \in c(|\{e \mid(e \in c(e b d))\}|)))$
D.8.a $=b \Leftrightarrow(\forall c, d($ string $[c, d](a \in d$
$D .15 . S_{b}^{a}(c)=d \Leftrightarrow(\operatorname{string}[a, b](c \in b(d \in b(\operatorname{cad})(-((\exists f(\operatorname{fad})(c a f)))$

```
D.16. \(a \simeq^{b, c} d \Leftrightarrow(\forall c, d(\) string \([c, d](a \in d(c, d((\neg(c, d \Rightarrow(g b h(g \in c(h \in c)))))\)
                                    ( \(c, d[a \mid b](\forall e, f(\) string \([e, f](\)
    \(b \in f(e, f((\neg(e, f \Rightarrow(i b j(i \in f(j \in f)))))(e, f[b \mid a]))))))))))))\)
D.17. \((a=\{b, c\}(\mid \notin c(\) string \([b, c]))) \Leftrightarrow(d \in a \Leftrightarrow(((d=\inf (b, c)) \vee(d=\sup (b, c))) \vee(\exists e\)
                                    \(\left.\left.\left(e \simeq^{b, c}=,\left(\exists f\left(f \simeq^{b, c}=,\left(S_{c}^{b}(e)=d\left(S_{c}^{b}(d)=f\right)\right)\right)\right)\right)\right)\right)\)
                    D.18.(string \([b, c](a \in c(a=b, c))) \Leftrightarrow(c=\{a\})\)
\(D .19 .(a c b=c, d(\operatorname{string}[c, d])) \Leftrightarrow(b \in d(a=c, e(\operatorname{string}[c, e](e=d /\{b\}(\forall f(f \in e(f c b)))))))\)
                    D.20.deci \(=\{0,1,2,3,4,5,6,7,8,9\}\)
            D.21. \(<\), deci \(=0<1<2<3<4<5<6<7<8<9\)
    \(D .22 . \mathbb{N}_{a, b}=\left\{a, c \mid\left(\right.\right.\) string \(\left.\left.[a, c]\left(\forall d\left(d \in c\left(d \simeq^{a, c} e(e \in b)\right)\right)\right)\right)\right\}([a, b])\)
\(D .23 . a: b \rightarrow c \Leftrightarrow(\forall d(d \in b(\exists e(e \in c(a(d)=e(\forall f(f \notin b(\forall g(a(f) \neq g)))))))))\)
            D.24.injection \([a, b, c] \Leftrightarrow(a: b \rightarrow c(a(d)=a(e) \Rightarrow d=e))\)
        \(D .25\). surjection \([a, b, c] \Leftrightarrow(a: b \rightarrow c(\forall d(d \in c(\exists e(e \in b(a(e)=d))))))\)
            D.26.bijection \([a, b, c] \Leftrightarrow(\) injection \([a, b, c]\) (surjection \([a, b, c]))\)
                D.27. \(|a|<|b| \Leftrightarrow(\exists c(\) injection \([c, a, b]))\)
                \(D .28 .|a|=|b| \Leftrightarrow(\exists c(\) bijection \([c, a, b]))\)
D.29.a, \(b<\mathbb{N}_{a, c} a, d \Leftrightarrow\left(a, b \in \mathbb{N}_{a, c}\left(a, d \in \mathbb{N}_{c, d}(\right.\right.\) string \([a, b](\) string \([a, d]((|b|<|d|) \oplus\)
    \(\left.\left.\left.\left.\left.\left(|b|=|d|\left(\exists e\left(\forall f\left(f<e\left(f \operatorname{th}(a, b)=f \operatorname{th}(a, d)\left(\operatorname{eth}(a, b)<_{\mathbf{N}_{a, c}} \operatorname{eth}(a, d)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\)
                    \(D .30 \mathbb{N}_{a, b}(c)=d \Leftrightarrow\left(c \in \mathbb{N}_{a, b}\left(\left|\left\{e \mid\left(e<_{\mathbb{N}_{a, b}} c\right)\right\}\right|=d\right)\right)\)
                \(D .31 . \mathbb{N}=\left\{a \mid\left(\exists b\left(\exists c\left(\operatorname{string}[b]\left(\mathbb{N}_{b}(c)=a\right)\right)\right)\right)\right\}\)
                \(D .32 . T M[a, b, c, d, e, f, g] \Leftrightarrow(c \in b(d=b / c(g \subseteq a(f \in a(\)
    \(\left.\left.\left.\left.\left.e:(a / g) \times b \times\{T M L, T M R\} \rightarrow a \times b \times\{T M L, T M R\}\left(|g|=\mathbb{N}_{<, \text {deci }}(2)\right)\right)\right)\right)\right)\right)\)
\(D .33 . a \preccurlyeq_{b, c} d \Leftrightarrow\left(\left(\operatorname{set}[a] \vee \operatorname{class}[a]\left((\operatorname{set}[d] \vee \operatorname{class}[d])\left(\forall e\left(e \in a\left(e \simeq^{b, c} f(f \in d)\right)\right)\right)\right)\right)\right.\)
D.34.compfunc \([a, b] \Leftrightarrow(\operatorname{string}[, a](\forall c(b(c)=d(e(c))(\operatorname{compfunc}[f, e](d=\sup (a)(f=a /\{d\}))))))\)
\(D .35 . D T M[a, b, c, d, e, f, g, h, i, j] \Leftrightarrow(, j=A, B, C(A \in g(B \in b(C \in \mathbb{N}(T M[a, b, c, d, e, f, g](\)
    string \([, i]\left(\right.\) string \([, j]\left(i \lesssim b\left(\right.\right.\) compfunc \([h, k]\left(k\left(f,, i, \mathbb{N}_{<, \text {deci } i}(0)\right)=, j(|h| \in \mathbb{N}(\right.\)
\((l \operatorname{th}(, h)(m, n, o)=p, q, r) \Leftrightarrow\left(\left(\left(\left(o=\mathbb{N}_{<, \text {deci }}(0)\left(e(m, \inf (n))=p, \mathbb{N}_{<, \text {deci }}(1), T M L\right.\right.\right.\right.\right.\)
\(\left.\left.\left(r=\mathbb{N}_{<, \operatorname{dec} i}(0)\left(\inf (q)=c\left(\forall s\left(r<_{\mathbb{N}} s\left(\operatorname{sth}(n)=S_{\mathbb{N}}^{\in}(s) \operatorname{th}(q)\right)\right)\right)\right)\right)\right)\right) \oplus(\operatorname{oth}(n)=\sup (n)\)
\(\left.\left.\left(e(m, \sup (n))=p, \operatorname{sth}(q), \operatorname{TMR}\left(S_{\mathbb{N}}^{<}(0)=r\left(\sup (q)=\operatorname{rth}(q)\left(r \operatorname{th}(q)=c\left(\forall t\left(t<_{\mathbb{N}} o(\operatorname{tth}(n)=\operatorname{tth}(q))\right)\right)\right)\right)\right)\right)\right)\right)\)
\(\oplus\left(o \neq \mathbb{N}_{<, \operatorname{dec} i}(0)(e(m, \operatorname{oth}(n))=p, \operatorname{oth}(q), T M L(\forall s(s \neq o(\operatorname{sth}(n)=\operatorname{sth}(q)))))\right) \oplus(\operatorname{oth}(n) \neq \sup (n)\)
    \((e(\operatorname{m}, \operatorname{oth}(n))=p, \operatorname{oth}(q), \operatorname{TMR}(\forall s(s \neq o(\operatorname{sth}(n)=\operatorname{sth}(q)))))))))))))))))))\)
            D.36.NDTM \([a, b, c, d, e, f, g, h, i, j] \Leftrightarrow(\forall k(k \in e(T M[a, b, c, d, k, f, g](\)
                    compfunc \([k, l]\left(l\left(f, l, \mathbb{N}_{<, \text {deci }}(0)\right)=, j\right.\)
\((, j=A, B, C(A \in g(B \in b(C \in \mathbb{N}(\forall m(m<\mathbb{N}|h|(\exists n)(n \in e(m t h(h)(o, p, q)=r, s, t\)
\(\left(\left(\left(n(o, \operatorname{qth}(p))=r, q \operatorname{th}(s), u\left(\left(q \neq \mathbb{N}_{<, \operatorname{dec} i}(0)\left(u=\operatorname{TML}\left(S_{\mathbb{N}}^{<}(t)=q\right)\right)\right) \oplus(q \operatorname{th}(p) \neq \sup (p)(u=T M R\right.\right.\right.\right.\)
\(\left.\left.\left.\left(S_{\mathrm{N}}^{<}(q)=t\right)\right)\right)\right) \oplus\left(n(o, q \operatorname{th}(p))=r, u \operatorname{th}(s), v\left(u=\mathbb{N}_{<, \operatorname{deci}}(1)\left(\inf (s)=c\left(\forall w\left(w \operatorname{th}(p)=S_{\mathrm{N}}^{<}(w) \operatorname{th}(s)(v=T M L\right.\right.\right.\right.\right.\)
\(\left.\left.\left.\left.\left.\left.\left(S_{\mathbf{N}}^{<}(t)=q\right)\right)\right)\right)\right)\right)\right) \oplus\left(u=q\left(S_{\mathbf{N}}^{<}(q) \operatorname{th}(s)=\sup (s)(q \operatorname{th}(p)=\sup (p)(v=\operatorname{TMR}(\sup (s)=c(\right.\right.\)
            \(\forall w(\boldsymbol{w}<\mathbf{N} q(\boldsymbol{w} \operatorname{th}(p)=w \operatorname{th}(s))))))))))))))))))))))))\)
D.37. \(\left(a=b^{c}(a \in \mathbb{N})\right) \Leftrightarrow\left(\right.\) compfunc \(\left.[d, e]\left(e\left(\mathbb{N}_{<, \text {deci }}(1)\right)=a(|d|=c(\forall f(\operatorname{fth}(, d)(g)=\mathrm{iN}(b, g))))\right)\right)\)
D.38.a(b) \(=O_{c}(d(b)) \Leftrightarrow\left(\exists e\left(\exists f\left(e \in \mathbb{N}\left(f \in \mathbb{N}\left(\forall g\left(g \in c\left(a(g) \leqslant \mathbb{N} \mathbb{N}\left(e, d(g)^{f}\right)\right)\right)\right)\right)\right)\right)\right)\)
                D.39. \(|a| \leqslant|b| \Leftrightarrow((|a|=|b|) \oplus(|a|<|b|))\)
            D.40. \(+_{\mathbb{N}}(a, b)=c \Leftrightarrow(c \in \mathbb{N}(a \in \mathbb{N}(b \in \mathbb{N}(\{d \mid(d<\mathbb{N} c(a \leqslant \mathbb{N} d))\}))))\)
            D.41. \({ }_{\text {iv }}(a, b)=c \Leftrightarrow(c \in \mathbb{N}(a \in h d s N(b \in \mathbb{N}\)
\(\left(\right.\) compfunc \(\left.\left.\left.\left.[d, e]\left(|d|=b\left(e\left(\mathbb{N}_{<, \operatorname{deci}}(0)\right)=c\left(\forall f\left(\forall g\left(f t h(d)(g)=+_{\mathbf{N}_{<, \text {deci }}}(a, g)\right)\right)\right)\right)\right)\right)\right)\right)\right)\)
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D.42.invath $(b, c)=d \Leftrightarrow(\mid\{e \mid$ dae $\} \mid=a$
$D .43 . P=\{k \mid((\operatorname{set}[k] \vee \operatorname{class}[k])(\exists a(\exists b)(\exists c(\exists d(\exists e(\exists f(\exists g(\exists l)(\exists m)(\forall i(\exists h(\exists j) N D T M[a, b, c, d, e, f, g, h, i, j])$

$$
\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\begin{array}{c}
|h|=\mathrm{iv}\left(l,|i|^{m}(j=n, o, p(n=\inf (, g) \Leftrightarrow\right. \\
i \notin k)((n=\sup (, g)
\end{array} \Leftrightarrow i \in k\right)(n=\sup (, g) \oplus n=\inf (, g))\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\}
$$

D.44.NP $=\{k \mid(\exists a(\exists b)(\exists c(\exists d) \exists e(\exists f(\exists g) \exists l(\exists m)(\forall i(, i \in k)(\exists h(\exists j) N D T M[a, b, c, d, e, f, g, h, i, j]$

$$
\left(|h|=\mathrm{in}^{\mathrm{N}}\left(l,|i|^{m}\right)(j=n, o, p(n=\inf (, g) \Leftrightarrow i \notin k)(\right.
$$

$$
(n=\sup (, g) \Leftrightarrow i \in k)(n=\sup (, g) \oplus n=\inf (, g))))))))))))))))))))\}
$$

$D .45 . \mathbb{Z}_{a, b}=\left\{a, c \mid\left(\operatorname{string}[a, c]\left(a, c \in \mathbb{N}_{a, b} \oplus\left(a, c=-a d\left(\operatorname{string}[a, d]\left(a, d \in \mathbb{N}_{a, c}\right)\right)\right)\right)\right)\right\}$

$$
\begin{gathered}
D .46 . \mathbb{Z}_{a}(b)=c \Leftrightarrow\left(\mathbb{N}_{a}(b)=c \oplus\left(b=-\operatorname{de}\left(\text { string }[d, e]\left(c=-\mathbb{N}_{a}(d, e)\right)\right)\right)\right) \\
D \cdot 47 . \mathbb{Z}=\left\{a \mid\left(\exists b\left(\exists c\left(\mathbb{Z}_{b}(c)=a\right)\right)\right)\right\}
\end{gathered}
$$

$D .48 \cdot \mathbb{R}_{a, b}=\left\{a, c \mid\left(\operatorname{string}[a, c]\left(\left(a, c \in\right.\right.\right.\right.$ mathds $Z_{a, b} \oplus\left(. \in c\left(\forall e\left(e \in c\left(\sup (a, c) \neq\left(\forall g\left(g \in c\left(g \simeq^{a, c} h(g \neq .(h \in b\right.\right.\right.\right.\right.\right.\right.$

$$
))))))))) \oplus\left(\inf (a, c)=-\left(\sup (a, c) \neq .\left(\forall g\left(g \in c\left(g \neq-\left(g \neq .\left(g \simeq^{a, c} h(h \in h(h \in b))\right)\right)\right)\right)\right)\right)\right)\right)\right)\right\} \Leftrightarrow \operatorname{string}[a, b]
$$

$$
D .49 \cdot \mathbb{R}_{a}(b)=\mathbb{Z}_{a}(b) \Leftrightarrow b \in \mathbb{Z}_{a}
$$

$D .50 \cdot \cdot_{\mathbf{Z}_{a, b}}^{+}(a, c, a, d)=a, e \Leftrightarrow\left({({\underset{\mathbb{Z}}{a, b}}}(a, c, a, d)=a, e \oplus\left({\underset{\mathbf{Z}_{a, b}}{+}}(a, c /\{\inf (a, c)\}, a, d /\{\inf (a, d)\})=a, e\right.\right.$
$\left.(\inf (a, c)=-(\inf (a, d)=-))) \oplus\left(\inf (a, d)=-\left(\inf (a, c) \neq-\left(-\mathbb{N}_{a, b}(a, c, a, d /\{\inf (a, d)\})\right)\right)\right)\right)$

$$
D .51 \cdot \mathbb{R}=\left\{a \mid\left(\exists d\left(\exists b, c\left(\text { string }[b, c]\left(\mathbb{R}_{b, c}(d)=a\right)\right)\right)\right)\right\}
$$

$D .52 \cdot \cdot_{\mathbf{R}_{a, b}}^{+}(c, d)=e \Leftrightarrow\left(\left(\left(\hat{+}_{\mathbf{Z}_{a, b}}(c, d)=e\right) \oplus(((\inf (c)=-(\inf (d)=-)) \oplus(\inf (c) \neq-(\inf (d) \neq-)))\right.\right.$
$\left((\neg(\exists f(\operatorname{fth}(c)=))).\left(g \operatorname{th}(d)=.\left(c=a, i\left(d=a, j\left(_{\mathbf{Z}_{a, b}}(c,\{l \mid(l=m t h(d)(m<N g))\})=n(\right.\right.\right.\right.\right.$
$\left.\left.\left.\left.\left.\left.\left.\left.e=a, n \cup\left\{m \mid\left(m=\operatorname{oth}(d)\left(g \leqslant_{\mathbb{N}} o\right)\right)\right\}\left(\sup (a, n) a\left(a,\left\{m \mid\left(m=\operatorname{oth}(d)\left(g \leqslant_{\mathbb{N}} o\right)\right)\right\}\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)$
$\oplus(((\inf (c)=-(\inf (d)=-)) \oplus(\inf (c) \neq-(\inf (d) \neq-)))(\operatorname{fth}(c)=.(g t h(d)=.(f \leqslant \mathbb{N} g(\sup (c)=j \operatorname{th}(c)$
$\left(\sup (d)=k t h(d)\left(\operatorname{compfunc}[l, m]\left(m(d)=e\left(n t h(, l)(d)=o \Leftrightarrow\left(\left(\left(n<_{\mathbf{N}}-\mathbb{N}(k, f)\right.\right.\right.\right.\right.\right.\right.$

$(\forall p(p \neq n(\operatorname{invpth}(o)=\operatorname{invpth}(d)))))) \oplus\left(n<_{\mathrm{N}}-_{\mathbf{N}}(k, g)\left(-_{\mathbf{N}}(j, f)<_{\mathrm{N}} n\right.\right.$
$\left(\operatorname{invnth}(o)=\stackrel{+}{\mathbf{N}_{a, b}}\left(\operatorname{inv} S_{\mathbf{N}}^{<} n t h(c), \operatorname{invnth}(d)\right)\left({\stackrel{+}{\mathbf{N}_{a, b}}}_{+}\left(\operatorname{inv} S_{\mathbf{N}}^{<} n t h(c), \operatorname{invnth}(d)\right)=a, q\right.\right.$
$\left.\left.\left.\left.\left.\left.\left(|q|=\mathbb{N}_{<, \operatorname{deci}}(1)(\forall p(p \neq n(\operatorname{invpth}(o)=\operatorname{invpth}(d))))\right)\right)\right)\right)\right)\right) \oplus(d=o(\operatorname{invnth}(d)=g \operatorname{th}(d)))\right) \oplus$
$\left(n<_{\mathbf{N}}-_{\mathbf{N}}(k, f)\left(\operatorname{invnth}(o)=\sup \left({\stackrel{+}{\mathbf{N}_{a, b}}}(\operatorname{invnth}(c), \operatorname{invnth}(d))\right)\left({\stackrel{+}{\mathbf{N}_{a, b}}}^{(\operatorname{invnth}(c), \operatorname{invnth}(d))}=a, q(\right.\right.\right.$
$|q|=\mathbb{N}_{<, \operatorname{deci}}(2)\left(\forall p\left(p \neq n\left(\operatorname{invpth}(o)=\operatorname{invpth}\left({\stackrel{+}{\mathbf{N}_{a, b}}}_{(\inf (a, q), d)))))))))}\right.\right.\right.\right.$
$\oplus\left(n<_{\mathbf{N}}-\mathbf{N}(k, g)\left(-\mathbf{N}(j, f)<_{\mathbf{N}} n\left(\operatorname{invnth}(o)=\sup \left({\stackrel{+}{\mathbf{N}_{a, b}}}^{+}\right.\right.\right.\right.$
$\left.\left(\operatorname{inv} S_{\mathbf{N}}^{<}(n) \operatorname{th}(c), \operatorname{invnth}(d)\right)\right)\left({\stackrel{+}{\mathbf{N}_{a, b}}}\left(\operatorname{inv} S_{\mathbf{N}}^{<}(n) \operatorname{th}(c), \operatorname{invnth}(d)\right)=a, q\left(|q|=\mathbb{N}_{<, \text {deci }}(2)\right.\right.$
$\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left.\left(\forall p\left(p \neq n\left(\operatorname{invpth}(o)=\operatorname{invpth}\left({\stackrel{+}{\mathbf{N}_{a, b}}}(\inf (a, q), d)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)\right)$
$D .53 \cdot{ }_{\cdot}^{+}(a, b)=c \Leftrightarrow\left(\exists d, e\left(\operatorname{string}[d, e]\left(\exists f\left(\exists g\left(\exists h\left({\stackrel{1}{\mathbf{R}_{d, e}}}^{+}(f, g)=h\left(\mathbb{R}_{d, e}(f)=a\left(\mathbb{R}_{d, e}(g)=b\left(\mathbb{R}_{d, e}(h)=c\right)\right)\right)\right)\right)\right)\right)\right)\right.$ $D .54 \cdot \cdot_{\mathbf{R}}^{+}(a, b)={ }_{\mathbf{R}}^{+}(b, a)$
D.55. $\mathbb{Z}_{a, b}(a, c, a, d)=a, e \Leftrightarrow\left(\left(\left(\mathrm{iv}_{a, b}(a, c, a, d)=a, e\right)\right.\right.$
$\left.\oplus\left(a, c \in \mathbb{N}_{a, b}\left(\inf f(a, d)=-\left(\operatorname{iN}_{a, b}(a, c, a, d /\{\inf (a, d)\})=a, f(a, e=-a a, e)\right)\right)\right)\right)$
$\left.\oplus\left(\inf (a, c)=-\left(\inf (a, d)=-\left(\mathbf{i n}_{a, b}(a, c /\{\inf (a, c)\}, a, d /\{\inf (a, d)\})\right)\right)\right)\right)$
$D .56 . \dot{\mathbf{Z}}(a, b)=c \leftrightarrow\left(\exists d\left(\exists e\left(\exists f\left(\exists g\left(\mathbb{Z}_{d}(e, f)=g\left(\mathbb{Z}_{d}(e)=a\left(\mathbb{Z}_{d}(f)=b\left(\mathbb{Z}_{d}(g)\right)\right)\right)\right)\right)\right)\right)\right)$
D.57./ $\mathbb{Z}_{a}(b, c)=d \Leftrightarrow\left(\dot{\mathbf{Z}}_{a}(d, c)=b\right)$
D.58./ $\mathbb{Z}(a, b)=c \Leftrightarrow(\dot{\mathbf{Z}}(c, b)=a)$
\left. D.59.( ${\underset{\mathbb{R}}{a, b}}(a, c, a, d)=a, e\left(a, c \in \mathbb{Z}_{a, b} \oplus a, d \in \mathbb{Z}_{a, b}\right)\right) \Leftrightarrow((h(i)=j) \Leftrightarrow(i \in b$
$\left.\left.\left(\operatorname{iasup}(a, b)\left(S_{\mathbf{R}_{a, b}}^{<\mathbf{R}_{a, b}}(\sup (a, b))=k\left({ }^{\mathbb{N}_{a, b}}(i, j)=k\right)\right)\right)\right)\right)$
$\left(\dot{\mathbf{z}}_{a, b}(a, c, a,\{l \mid(l=\operatorname{mth}(a, d)(\boldsymbol{m}<g))\})=l(\operatorname{compfunc}[m, n]\right.$
$\left.\left.\left.\left.m(l))=a, e\left(\left(\operatorname{qth}(, m)(r)=s \Leftrightarrow s=/ \mathbf{z}_{a, b}\left(r, h\left(\mathbf{N V}^{+}(q, f) \operatorname{th}(a, d)\right)\right)\right)\right)\right)\right)\right)\right)$
$\left.D .60 .\left(\mathfrak{R}_{a, b}(a, c, a, d)=a, e\left(\neg\left(a, c \in \mathbb{Z}_{a, b} \oplus a, d \in \mathbb{Z}_{a, b}\right)\right)\right)\right) \Leftrightarrow\left(\left(\left(\dot{\mathbf{Z}}_{a, b}(a, c, a, d)=a, e\right) \oplus\right.\right.$
$(\operatorname{fth}(a, c)=.(\operatorname{gth}(a, d)=.(((h(i)=j) \Leftrightarrow(i \in b$
$\left.\left.\left(\operatorname{iasup}(a, b)\left(S_{\mathbf{R}_{a, b}}^{<\mathbf{R}_{a, b}}(\sup (a, b))=k\left({ }^{\mathbb{N}_{a, b}}(i, j)=k\right)\right)\right)\right)\right)$
$\left(\overleftarrow{\mathbf{Z}}_{a, b}(a,\{l \mid(l=\operatorname{mth}(a, c)(m<f))\}, a,\{l \mid(l=\operatorname{mth}(a, d)(m<g))\})=l(\right.$ compfunc $[m, n]($ compfunc $[o, p]$

```
(o(m(l)) =a,e((qth(,m)(r)=s\Leftrightarrows=/\mp@subsup{\mathbf{z}}{a,b}{}(r,h(\stackrel{+}{\mathbf{N}}=(q,f)\operatorname{th}(a,d)))(qth(,o)(r)=s
\Leftrightarrows=/\mp@subsup{\mathbf{z}}{a,b}{}(r,h(\mathbf{N}=(q,g)th(a,c)))))))))))))))
    D.61.BPP = {a| (\existsx (\existsy (x\in\mathbb{R}(y\in\mathbb{R}(\existsb,c,d,e,f,g,h, (\foralli(j(i)=|{l|)}\exists\textrm{\existsm}
        NDTM[b,c,d,e,f,g,h,l,k,m](|l| \leqslant <\mathbb{R R}
(((k\ina) => (\dot{\mathbb{R}}
(m=n,o,p(n=\operatorname{sup}(,h)(|l|\mp@subsup{\leqslant}{\mathbf{R}\mathbf{R}}{\mathbf{R}}(x,|k\mp@subsup{|}{}{y}))))))}|)))((k\not\ina(k\in\mp@subsup{c}{}{*}))}=>(|{l|(\existsm(NDTM[b,c,d,e,f,g,h,l,k,m
(m=n,o,p(n=\operatorname{sup}(,h)(|l|\leqslant\mathbb{R R}
D.62.a,bDTM[c,d,e,f,g,h,i]\Leftrightarrow(\forallj(\existsk(\existsl(DTM[c,d,e,f,g,h,i,k,j,l](|k|\leqslantb(|j| (l=m,n,o
                        ((m}=\operatorname{sup}(,i)\Leftrightarrowj\ina
(m=inf(,i)\Leftrightarrow(j\not\ina))))))))))
D.63.(ab
    D.64.a \in polyfunc}\Leftrightarrow(\existsd(\existse(\forallb(a(b)=c\Leftrightarrow(c=\textrm{iv}(d,\mp@subsup{b}{}{e})))))
    Ax.1.(\existsa(set[a]))
    Ax.2.(\existsa(class[a]))
Ax.4.(\existsa(a= IN <,deci})
Ax.5.(\existsa(a=\mathbb{N}))
    Th.1.P\subseteqBPP
    Pr.1.((a\inP) => (\existsb,c,d,e,f,g,h(a,jDTM[b,c,d,e,f,g,h](j\inpolyfunc))))
```



```
        j(i)=|{l|(\existsm(NDTM[b,c,d,e,f,g,h,l,k,m](|l|\mp@subsup{\leqslant}{\textrm{R}}{\textrm{R}}\mathbf{R}(x,|k\mp@subsup{|}{}{y}))))}|
    (((k\ina) => (% ( 
    (m=n,o,p(n=\operatorname{sup}(,h)(|l|\mp@subsup{\leqslant}{\mathbb{R}\mathbf{R}}{}(x,|k\mp@subsup{|}{}{y}))))))}|)))((k\not\ina(k\in\mp@subsup{c}{}{*}))=>(|{l|(\existsm(NDTM[b,c,d,e,f,g,h,l,k,m]
    (m=n,o,p(n=\operatorname{sup}(,h)(|l|\leqslant\mp@subsup{\leqslant}{\mathbb{R}}{\textrm{R}}
Th.2.(\existsa(a\subseteqP(string[,a](string [,b] ( |a| = |b| (\forallj(\forallk(\existsl(DTM[c,d,e,f,g,h,i,jth(,b),k,l]
(l=m,n,o((m=\operatorname{sup}(,i)\Leftrightarrow(k\injth(,a)))((m=\operatorname{inf}(,i)\Leftrightarrow(k\not\injth(,a)))(\forallj(NDTM[c,d,e,f,b,h,i,j,k,l]
(|j|\leqslant\mp@code{R R }}\mp@subsup{\mathbb{R}}{<,\mathrm{ deci}}{}(10),|k\mp@subsup{|}{}{\mp@subsup{\mathbb{R}}{<,deci}{\prime}(100)})))({k|(k\ina\vee(NDTM[c,d,e,f,b,h,i,j,k,l
(l=m,n,o(m=\operatorname{sup}(,i)(\existsA(\existsB(B\in\operatorname{polyfunc}(A,BNDTM[c,d,e,f,b,h,i,j,k,l]
(\neg(\existsC(C\in\operatorname{polyfunc(\existsD(A,CDTM[c,d,e,f,Dth(,)b,h,i,j,k,l])))))))))))))))}\not\inP}\\mp@code{)})))))))))))
Pr.2. ((()((a(b,c) =d\Leftrightarrowbth(,c) =d) (a, = {a(z,b)|(z\in\mathbb{N}(b\inP(b,dDTM[e,f,g,h,yth(,i),j,k]
```



```
o,pNDTM[e,f,g,h,q,j,k](\forallr(r\ino(r\ina(\existsA(\existsB(B\in\operatorname{poly func}(A,BNDTM[c,d,e,f,b,h,i,j,k,l]
(\neg(\existsC(C\in\operatorname{polyfunc}(\existsD(A,CDTM[c,d,e,f,Dth(,b),h,i,j,k,l])))))))))))))))))))))))))))}({01}\in\mp@subsup{a}{1}{}(\mp@subsup{a}{1}{}\inP
    ))))}=>(a\subseta))
((f1}(d,0)=0,0,TMR(\mp@subsup{f}{1}{}(0,1)=\operatorname{sup}(,h),1,TMR(\mp@subsup{f}{2}{}(d,1,)=0,0,TMR(\mp@subsup{f}{2}{}(0,0)=\operatorname{sup}(,h),0,TM
(f
f
f
    ))))})
b,dDTM[{0,1,2},{0,1},g,h,i,j,k]((i=\mp@subsup{f}{1}{}\oplusi=\mp@subsup{f}{2}{})(d\in\operatorname{polyfunc}(\foralll(l\in\mp@subsup{a}{1}{}(\existsm(m\in\mathrm{ poly func(})
```



```
(\forallr(r\ino(r\in\mp@subsup{a}{1}{}))))))))))))))))}))}({01}\in\mp@subsup{a}{1}{}))))))\wedge((a(b,c)=d\Leftrightarrowbth(,c)=d)
\existsa, (a, = {a(z,b)|(z\in\mathbb{N}(b\inP(b,dDTM[{0,1,2},{0,1},g,h,i,j,k]((i=\mp@subsup{f}{1}{}\oplusi=\mp@subsup{f}{2}{})(d\in\operatorname{polyfunc}(\foralll(l\in\mp@subsup{a}{1}{}(\existsm
m\in polyfunc( }\existsn(n\inq(l,mDTM[{0,1,2},{0,1},g,h,n,j,k
    (\forallo(\forallp(p\in polyfunc(o,pNDTM[{0,1,2},{0,1},g,h,q,j,k](\forallr(r\ino(r\ina(
```

```
    \existsA(\existsB(B) polyfunc(A,BNDTM[c,d,e,f,b,h,i,j,k,l]
(\neg(\existsC(C\in\operatorname{polyfunc}(\existsD(A,CDTM[c,d,e,f,D\operatorname{Dh}(,b),h,i,j,k,l]))))))))))))))))))))))))))))))}({01}\in\mp@subsup{a}{1}{})))}
((\neg)(((a)b,c) =d\Leftrightarrowbth(,c)=d) (a, = {a(z,b)|(z\in\mathbb{N}(b\inP(b\subseteqy(b,dDTM[e,f,g,h,i,j,k](d\in polyfunc
(\foralll(l\in\mp@subsup{a}{1}{}(\existsm)}(m\in\operatorname{polyfunc}(\existsn)(n\inq(l,mDTM[e,f,g,h,n,j,k])\forallo(\forallp(p\in\mathrm{ polyfunc(
```



```
    (\neg(\existsC)C 的oly func}(\existsD(A,CDTM[c,d,e,f,Dth(,b),h,i,j,k,l
```



```
((a(b,c) =d\Leftrightarrowbth(,c) =d) (\foralla, (a, = {a(z,b)|(z\in\mathbb{N}(b\inP(b,dDTM[e,f,g,h,i,j,k](d\in polyfunc
(\foralll(l\ina(\existsm)
```



```
(-(\existsC(C\in\operatorname{polyfunc}(\existsD(A,CDTM[c,d,e,f,Dth(,b),h,i,j,k,l]))))))))))))))))))))))))))))}({01}\in\mp@subsup{a}{1}{}
                                    (a
\wedge((\neg)((a(b,c) =d\Leftrightarrowbth(,c) =d) (a, = {a(z,b)|(z\in\mathbb{N}(b\inP(b\subseteqy(b,dDTM[e,f,g,h,i,j,k](d\in polyfunc
(\foralll(l\in\mp@subsup{a}{1}{}(\existsm)
```



```
    (\neg(\existsC)C\in\mathrm{ poly func }(\existsD(A,CDTM[c,d,e,f,Dth(,b),h,i,j,k,l]
```



```
((a(b,c) =d \Leftrightarrowbth(,c) =d) (\forall\mp@subsup{a}{1}{}(\mp@subsup{a}{1}{}={a(z,b)|(z\in\mathbb{N}(b\inP(b,dDTM[e,f,g,h,i,j,k](d\in\mathrm{ polyfunc}
(\foralll(l\ina(\existsm)
```



```
    (}\neg(\existsC(C\in\mathrm{ poly func }(\existsD(A,CDTM[c,d,e,f,Dth(,b),h,i,j,k,l
```



```
        Th.3.(\existsa(a\subseteqP(a\not\inP(a\inBPP))))
Pr.3(a\subseteqP(string[,a](string[,b] (|a| = |b| (\forallj) \forallk(\existsl(DTM [c,d,e,f,g,h,i,jth(,b),k,l]
(l=m,n,o((m=\operatorname{sup}(,i)\Leftrightarrow(k\injth(,a)))((m=\operatorname{inf}(,i)\Leftrightarrow(k\not\injth(,a)))(\forallj(NDTM[c,d,e,f,b,h,i,j,k,l]
(|j| \leqslant {\mathbb{R}\mathbb{R}
(l=m,n,o(m=\operatorname{sup}(,i)(\existsA(\existsB(B\in\operatorname{polyfunc}(A,BNDTM[c,d,e,f,b,h,i,j,k,l](\neg)(\existsC)C\in\operatorname{polyfunc}(\existsD)
A,CDTM[c,d,e,f,\operatorname{Dth}(,b),h,i,j,k,l])))))))))))))))}=z(z\not\inP())))))))))))))=>(z\inBPP))\wedge
                    (\existsa(a\subseteqP(string [,a](string[,b]
    ( |a| = |b| (\forallj(\forallk(\existsl(DTM[c,d,e,f,g,h,i,jth(,b),k,l]
(l=m,n,o((m=sup(,i)\Leftrightarrow(k\injth(,a)))((m=\operatorname{inf}(,i)\Leftrightarrow(k\not\injth(,a)))(\forallj(NDTM[c,d,e,f,b,h,i,j,k,l]
(|j|\leqslant\mp@code{R R }
(l=m,n,o(m=\operatorname{sup}(,i)(\existsA(\existsB(B\in\operatorname{polyfunc}(A,BNDTM[c,d,e,f,b,h,i,j,k,l]
                            (})(\existsC(C\in\mathrm{ polyfunc }(\existsD
A,CDTM[c,d,e,f,\operatorname{Dth}(,b),h,i,j,k,l])))))))))))))))}=z(z\not=P(\mp@code{))}))))))))))))
                                    Th.4.P}\subsetBP
    Pr.4.((\existsa(a\not\inP(a\inBPP))) ^(P\subseteqBPP))
```


## VI. English nomenclature

$a \mathrm{~L} b$ in this section means every formal mathematical statement in the languageof the paper about $b$ has an equivalent formation in the english language about

## N. 1

a is an subset of $\mathrm{b} \Leftrightarrow a \subseteq b$
N. 2
a iff $\mathrm{b} \Leftrightarrow(a \Leftrightarrow b)$
a is a set $\Leftrightarrow \operatorname{set}[a]$
N. 3
N.4.
a is a state set for some deterministic configuration that solves b

```
            \Leftrightarrow(\existsc(\existsd)}\exists\textrm{e}(\exists\textrm{f}(\existsg)\existsh(\foralli(\existsj)(\existsk)DTM[a,c,d,e,f,g,h,j,i,k
(k=l,m,n((n=inf(,h)\oplusn=\operatorname{sup}(,h))((n=inf(,h)\Leftrightarrow,i\not\inb)(n=\operatorname{sup}(,h)\Leftrightarrow,i\inb))))))))))))))
                                    N.5.
a is a state set for some non deterministic configuration that solves b
            \Leftrightarrow(\existsc(\existsd)(\existse(\existsf)}\exists\textrm{g}(\exists\textrm{h}(\foralli)(\existsj)\existsk(NDTM[a,c,d,e,f,g,h,j,i,k
(k=l,m,n((n=\operatorname{inf}(,h)\oplusn=\operatorname{sup}(,h))((n=\operatorname{inf}(,h)\Leftrightarrow,i\not\inb)(n=\operatorname{sup}(,h)\Leftrightarrow,i\inb))))))))))))))
    N.6.
a is an element of }\textrm{b}\Leftrightarrowa\in
    N.7
there exists a such that b}\Leftrightarrow(\existsa(b)
    N.8.
for every a such that b c}\Leftrightarrow(\foralla(b(c))
    N.9.
if a then b
    N.10.
It is not true that a}\Leftrightarrow(\square(a)
    N.11.
a is not an element of b}\Leftrightarrowa\not\in
    N.12.
a}\mathrm{ is a superset of b}\Leftrightarrowb\subseteq
cardinality of a L |a|
    N.14.
delta 2 of a that solves b L }\mp@subsup{\delta}{b}{2}(a
    N.15.
a is a proper subset of b}\Leftrightarrowa\subset
    N.16
a is an improper subset of b
```


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