

Data Digest-based Authentication for Mobile Cloud Computing

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Abstract: Mobile cloud computing can influence the long run of various applications, like electronic commerce and health IP. With mobile cloud computing, resource-constrained mobile devices may maximize the computation/storage resources of cloud servers via communication networks, mobile devices in mobile cloud computing are having lot of security risks as a result they usually got to access cloud servers through untrusted networks from completely different locations. one among the foremost necessary aspects of mobile cloud computing security is to determine documented communication sessions between mobile devices and cloud servers., DDA strategically incorporates hashing, additionally to ancient user ID and passwords, to realize mutual authentication. The effectiveness of DDA is valid with Scyther, security protocol analyser. Our experimental results indicate that DDA is capable of withstanding a range of various-security-attacks.
Keywords: Cloud computing, Security, Mobile devices, Authentication, Hashing

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I. Introduction

In this paper, we have a tendency to commit to improve the safety of mobile cloud computing by introducing a unique authentication theme, Data Digest-based Authentication, a mobile cloud uses either user id and watchword primarily based authentication, or USIM (universal subscriber identity module) primarily based authentication [2]. DDA employs hashing and therefore the ancient user id/password to make a good authentication mechanism. DDA ensures that mobile cloud computing is secured from any kind of unauthorized access at the start of every communication session. DDA employs the shopper registration termination and unidirectional mathematical hashing to cypher the user ID so as to form the authentication method safer and a lot of unpredictable. Remote location of resources and virtualization technologies build the cloud computing settings susceptible to attacks. In cloud computing, all shoppers access a standard resource location that introduces security threat to the system. There's an associate integrity issue in cases of transfer, storage and retrieval. There is no common place to make sure knowledge integrity. There is an associate inherent vulnerability within the service offered. If the virtualization platform is compromised, it implies that a majority of the virtual machines are vulnerable, that could be a potential threat to knowledge security. To address the top mentioned problems, it's essential to possess a customary cloud computing style and usage policy, worker trust, proprietary computer code for virtualization. In a mobile cloud computing setting, a mobile device has got to be registered with a cloud server as a pre-requisite method before avail any types of cloud services. The information transmission between the mobile device and therefore the cloud server should be performed once the mobile device authenticates the cloud server and vice-versa. a powerful authentication theme ensures secure communication between 2 legitimate parties although the channel experiences potential vulnerability. Since the mobile device lacks of procedure capability, it's not appropriate to use complicated operations within the mobile device for authentication method. In mobile cloud computing setting, if a mobile device is registered with a selected cloud service supplier, each mobile device and cloud server should evidence one another in a very uniform manner so as to secure the communication with one authentication theme, that permits a mobile user to access the cloud server from totally completely different completely different locations mistreatment different networks and differing kinds of mobile devices.

H a: In mobile cloud computing setting, the projected authentication theme introduces hashing primarily based security, that reduces the vulnerability of the system to attacks. to see the vulnerability of the system, we have a tendency to work out the vulnerability score, S_v . the worth of S_v lies between zero.0 and 1.0:

$$H_a : 0. \text{zero} \leq S_v \leq 1.0 \quad (1)$$

A low score of S_v indicates that a theme offers a lot of security

Registration

The registration method of a mobile device or mobile user to a cloud server could be only one method whereby the user ID and therefore the watchword is setup and few encrypted files are changed. Upon fitting the mobile user account with cloud server, the cloud server performs a series of operations. Algorithm 1 shows the elaborated registration method adopted in our analysis.

Algorithm 1 Registration: mobile_with_cloud

Require:
isAlive (*mobile, cloud*)
hasNetworkAccess (*mobile*)

Ensure:
Role_Mobile
const *userID*
var *password*
uid = **hash**(*userID*)
pwd = **hash**(*password*)
cloud ← **send**(*uid*||*pwd*)
 $T_k = uid \oplus pwd$

Role_Cloud
recv(*uid*||*pwd*)
 $T_k = uid \oplus pwd$
EXP = Registration expiry period for the client
#CF ← **pointer**{**store**(*uid, pwd, T_k*)}
MD_{user} = **hash**(usage policy, user access level, user certificate)
MD_{cloud} = **hash**(user add policy, cloud resource restriction, cloud certificate)
MD = **hash**(*MD_{cloud}*||*MD_{user}*)
Temp = *MD_{user}*||*MD_{cloud}*||**#CF**||*Pk_{pub_cloud}*
msz = **encrypt**(*Temp, T_k*)
mobile ← **send**(*msz*)

The two parties concerned within the registration method are unit mobile device and the cloud server. Each of the parties should be alive or be within the network to accomplish the registration method. Cloud Server stores hash, hash, and user's mobile device info in massive table for economical operation. It generates 2 hashed messages or data digests. The first, *DD_{user}*, that consists of user policy (cloud resource usage policy, and user access level), and User certificate. The second message digest is *DD_{cloud}*. Upon generating each message digests, cloud server creates associate encrypted message to transmit these info to the mobile device. *ET_k*, where *Pk_{pub_cloud}* is that the cloud's public key, *DD_{user}* and *DD_{cloud}* area unit the generated message digests, and **#CF** is that the column reference, that refers to the cloud authentication information for that individual cloud user info. These info area unit sent from the cloud server to the mobile device when encrypting with key *T_k* that's generated in each the mobile device and cloud server by XOR-ing (Exclusive OR) hashed user ID and hashed parole (Eq. 2).

$$T_k = \text{hash} \oplus \text{hash} \tag{2}$$

Authentication

The cloud user has 2 Data digests *DD_{user}*, and *DD_{cloud}* within the mobile device. The authentication method is split into 2 steps: cloud authenticating mobile device and mobile device authenticating cloud.

3.1. Cloud authenticating mobile

When a mobile device desires to send associate authentication request to the cloud server, it generates a key *T_k*, victimization hashed user ID and hashed parole (Eq. 2). The key *T_k*, works because the seed for the PRNG (pseudo random range generator) to get associate authentication key, *Auth_Key_i*. This authentication key *Auth_Key_i* is needed to code the message digest *DD*, that is generated by hashing *DD_{cloud}* and *DD_{user}* (Eq. 3). The *Auth_Key_i* is that then the sequence of bits generated by PRNG that's such as by the state symbol *SI*. Key *T_k* is employed to code the state symbol *SI*, and therefore the encrypted message digest *EAuth_Key_i*. Finally, the encrypted message, *ET_k*{*EAuth_Key_i*||*SI*} is shipped to the cloud server (Fig. 1) along side the column reference **#CF**. Therefore, the message sent from mobile device to cloud server is **#CF**||*ET_k*{*EAuth_Key_i*{*MD*}||*SI*}
 $MD = \text{hash}\{MD_{cloud} || MD_{user}\}$

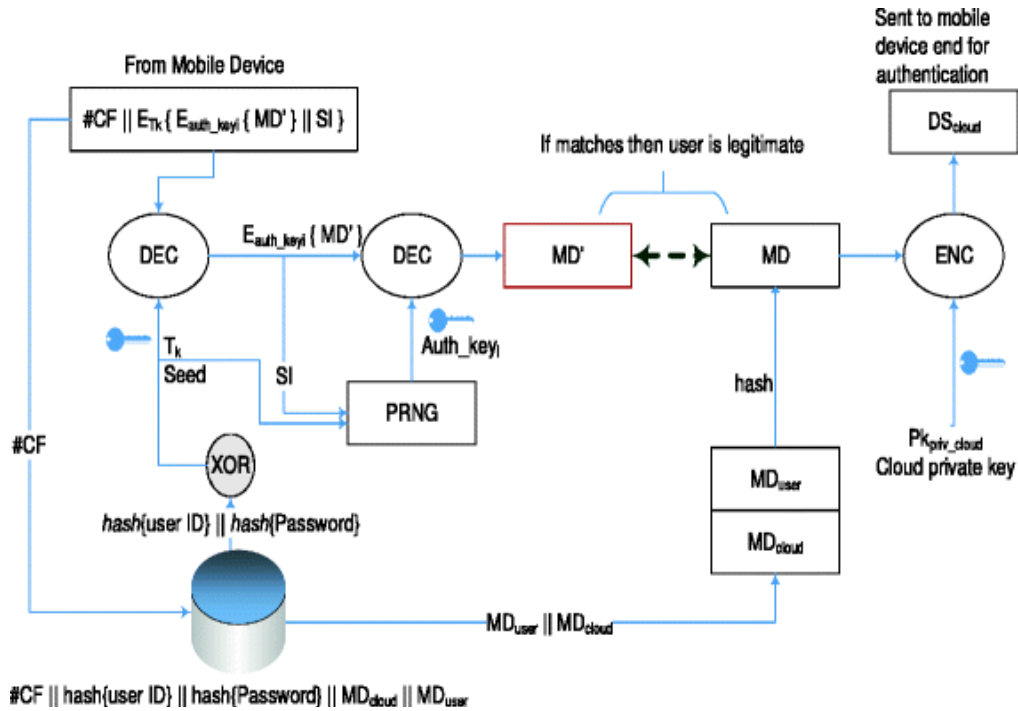


Fig. 1

Algorithm 2 Mobile_to_cloud_auth_req

Require:

isAlive (*mobile, cloud*)
hasNetworkAccess (*mobile*)
isRegistered (*mobile, cloud*)
 $uid = \mathbf{hash}(userID)$
 $pwd = \mathbf{hash}(password)$
 $cloud \leftarrow \mathbf{send}(uid \parallel pwd)$
 $T_k = uid \oplus pwd$

Ensure:

$MD = \mathbf{hash}(MD_{cloud} \parallel MD_{user})$
 $Auth_Key_i \leftarrow T_k \text{ PRNG } SI$
 $Enc_1 = \mathbf{encrypt}(MD, Auth_Key_i)$
 $Temp = Enc_1 \parallel SI$
 $Enc_2 = \mathbf{encrypt}(Temp, T_k)$
 $cloud \leftarrow \mathbf{send}(Enc_2 \parallel \#CF)$

Algorithm 3 Cloud_authenticating_mobile**Require:**

isAlive (*mobile, cloud*)
hasNetworkAccess (*mobile*)
isRegistered (*mobile, cloud*)

Ensure:

```

cloud ← recv(Enc2||#CF)
SearchEntity ← (uid||pwd)
search(SearchEntity, #CF)
if SearchEntity == found
    if EXP == expired
        triggerUpdatePhase()
    else
        Tk = uid ⊕ pwd
        Temp = decrypt(Enc2, Tk)
        Temp = Enc1||SI
        Auth_Keyi ← Tk PRNG SI
        MD = decrypt(Enc1, Auth_Keyi)
        MD' = hash(MDcloud||MDuser)
        if MD == MD'
            integrityCheck(pass)
            authentication(pass)
        else
            integrityCheck(fail)
            authentication(fail)
        endif
    endif
endif
else
    authentication(fail)
endif

```

Upon receiving the authentication request message, the cloud server performs decipherment operations. The cloud server searches the precise hashed userID and hashed parole within the cloud authentication information supported the shared column reference # C F that's sent in plain text together with the encrypted message. Once the hashed userID and hashed parole area unit found, the cloud server checks the registration validity of the user mistreatment EXP, keep at the server. If the user registration isn't invalid, then T_k is generated (Eq. 2) by XOR operation. The generated T_k at the cloud server decrypts the message $E_{Tk}||SI$, The authentication key, $Auth_Key_i$ decrypts the encrypted message digests $E_{Auth_Key_i}$. Algorithm 3 summarizes the operations concerned within the second sub-step.

3.2. Mobile authenticating cloud

Once the mobile device is documented, the cloud server sends its digital signature, that consists of DD encrypted with cloud's non-public key Pk_{priv_cloud} , to the mobile device. Alg 4 summarizes the operations performed at the cloud server throughout this section.

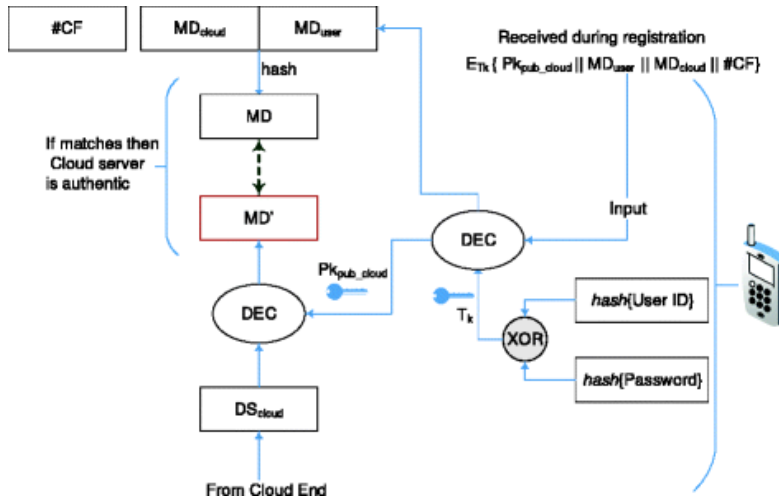


Fig. 2

Algorithm 4 Cloud_to_mobile_auth_req

Require:

- isAlive** (*mobile, cloud*)
- hasNetworkAccess** (*mobile*)
- isRegistered** (*mobile, cloud*)

Ensure:

- $MD = \text{hash}(MD_{cloud} || MD_{user})$
- $DS = \text{encrypt}(MD, Pk_{priv_cloud})$
- $mobile \leftarrow \text{send}(DS)$

After receiving the digital signature DS, the mobile device decrypts it with cloud’s public key Pk_{pub_cloud}, keep within the mobile device. If the decrypted DD matches with the message digest MD’, keep within the mobile device, then it may be declared that the cloud server is legitimate. Alg 5 provides the operations that performed by the mobile device once receiving cloud’s digital signature DS.

Algorithm 5 Mobile_authenticating_cloud

Require:

- isAlive** (*mobile, cloud*)
- hasNetworkAccess** (*mobile*)
- isRegistered** (*mobile, cloud*)
- $uid = \text{hash}(userID)$
- $pwd = \text{hash}(password)$
- $T_k = uid \oplus pwd$

- $MD = MD_{user} || MD_{cloud}$
- $msz \leftarrow \text{decrypt}((MD || \#CF || Pk_{pub_cloud}), T_k)$
- $msz = MD || \#CF || Pk_{pub_cloud}$

Ensure:

- $mobile \leftarrow \text{rcv}(DS)$
- $MD' = \text{decrypt}(DS, Pk_{pub_cloud})$
- if** $MD == MD'$
 - authentication**(*pass*)
 - integrityCheck**(*pass*)
- else**
 - authentication**(*fail*)
 - integrityCheck**(*fail*)
- endif**

Update

During the registration section the cloud server generates termination period for every mobile consumer. This termination period EXP is checked at the cloud server for every authentication request created by a mobile consumer. The new hashed number is going to be completely different than the recent hashed number and thereby the generated Tk (Eq. 2) are going to be completely different.

Evaluation-Methodology

We assume that the mobile device is already registered with the cloud server and obtained DD cloud, DD user, # C F, and Pkpub_cloud once a registered mobile sends the authentication request $\{SI\}\{CF\}$ to the cloud server, or once the cloud server sends authentication response Pkpriv_cloud to the mobile device, Methodology N is the variety of attacks that are unit launched on the planned theme and N success is that the variety of victorious attacks that are unit recorded. Then, the probability of victorious attacks on the theme defines its vulnerability score (Sv) (Eq. 4). Lesser vulnerability score indicates that the planned theme will stop a lot of variety of attacks, the value of Sv lies between zero and 1.0 (Eq. 1) $Sv = \frac{N_{success}}{N}$ We designed the Teddy boy setup and therefore the protocol analyser Scyther for collateral the DD theme. Scyther is designed from the setting choice to launch every type of attacks. Scyther tried to launch different kinds of attacks considering hackers have initial information of the system. The planned theme uses public key cryptography within the registration method, parallel cryptography in authentication method and a typical hashing algorithmic rule in generating the message digests. We have used RSA (Rivest, Shamir, and Adelman) the public key cryptography algorithmic rule, SHA1 (Secure Hash algorithmic rule 1) the hashing algorithmic rule, and AES (Advanced cryptography Standard) the parallel key cryptography algorithmic rule. The man-in-the-middle attack object is activated willy-nilly, and it changed the request frames that are unit sent. Among one thousand frames, 256 frames are unit changed in transit. The cloud server running DD theme has rejected all the changed request frames. mobile device, reset of EXP forces the cloud server to reject the authentication request. If EXP resets, a registration object is initialized to perform the re-registration.

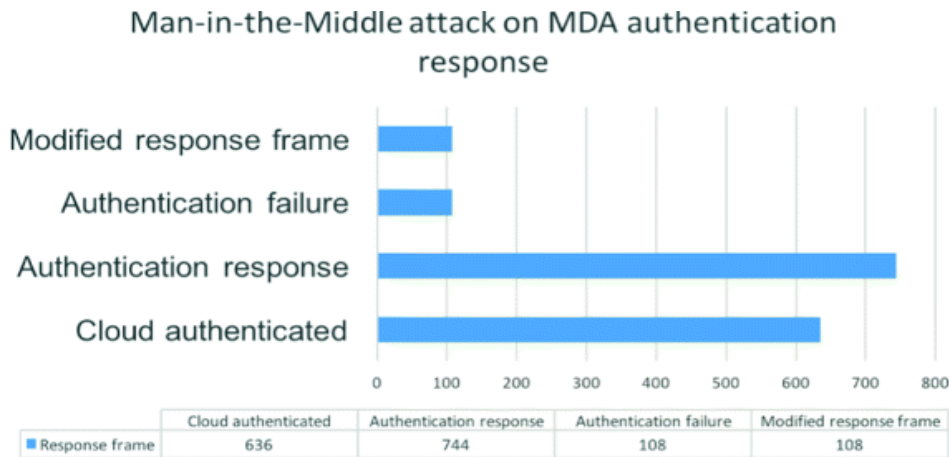


Fig. 3 Man-in-the-middle attack on DDA authentication response

Security Analysis Results

Table 1 summarizes our claims and how secure the scheme is for each claim. In this section, we provide detailed security analysis of the claims that are validated using Scyther.

Table 1 Security Analysis

| Role | Sl.Num. | Claim | Status | Comment |
|--------|---------|----------------------|--------|------------|
| Mobile | 1 | secret $Auth_Key_i$ | Ok | No Attacks |
| | 2 | secret SI | Ok | No Attacks |
| | 3 | secret $password$ | Ok | No Attacks |
| | 4 | secret $userID$ | Ok | No Attacks |
| | 5 | secret MD_{user} | Ok | No Attacks |
| | 6 | secret MD_{cloud} | Ok | No Attacks |
| | 7 | Alive | Ok | No Attacks |
| | 8 | Weak agree | Ok | No Attacks |
| | 9 | Ni agree | Ok | No Attacks |
| | 10 | Ni synchron | Ok | No Attacks |

| | | | | |
|-------|---|------------------------------------|----|------------|
| Cloud | 1 | secret <i>Auth_Key_i</i> | Ok | No Attacks |
| | 2 | secret <i>password</i> | Ok | No Attacks |
| | 3 | secret <i>userID</i> | Ok | No Attacks |
| | 4 | secret <i>SI</i> | Ok | No Attacks |
| | 5 | Alive | Ok | No Attacks |
| | 6 | Weakagree | Ok | No Attacks |
| | 7 | Niagree | Ok | No Attacks |
| | 8 | Nisynch | Ok | No Attacks |

Cond 1: *Auth_Key_i* remains secret throughout the authentication method.

Auth_Key_i is employed to write and decode the Data digest DD to produce multi layer security. A Mobile device sends the Data digest DD to the cloud server by encrypting with *Auth_Key_i*, that could be a radically symmetrical key. Each party, the mobile device and also the cloud server will severally generate the *Auth_Key_i* victimization state symbol *SI*. Since *Auth_Key_i* isn't changed between each of the parties, it remains secret.

Cond2: State symbol *SI* is secret. *SI* is that the state symbol for PRNG specifies the *n*th sequence of PRNG because of the desired pattern. The mobile device sends *SI* to the cloud server for specifying the *n*th sequence of PRNG to come up with the *Auth_Key_i*. *SI* is shipped from mobile device to cloud server when encrypting with key *Tk*. The claim at the mobile being secret is valid victimization Scyther. On the opposite hand, the cloud server doesn't share or send *SI*, therefore, it's safe at the cloud end.

Cond3: Password is secret. The safety of the secret password depends on the user. Generally nobody shares their secret, thereby keeping it safe. The secret password may be a 512 bit string, that is hashed and a duplicate of the user's secret is kept at the cloud server throughout registration method. The mobile device doesn't send the secret throughout authentication method, however it is used at each end to come up with the key *Tk*. Our secret password safety claim is valid by Scyther, but, in reality, it's shooked in to the user.

Cond 4: User Identification is confidential. *UserID* is shipped to the cloud server solely throughout the registration method, that is needed beside the secret to come up with *Tk*. The *UserID* is hashed and a duplicate is kept at the cloud server throughout registration method. The mobile device doesn't send *UserID* to the cloud server throughout authentication method. Therefore, *UserID* remains safe. Our *UserID* safety claim is valid by Scyther, but, in reality, it's shooked in to the user.

Cond 5: The theme needs DD user to be a secret. DD user is that the hashed data associated with the user, which can contain policy data and distinctive data regarding the user. This data is generated at the cloud server and sent to the mobile device when user sent registration method. This is hashed and encrypted with *Auth_key_i* before transmission from the user mobile device to the cloud finish. Scyther valid our claim that DD user is safe..

Cond 6: Mobile device and also the cloud server remains alive throughout the execution of the protocol. The cloud server is alleged to be alive if it's been victimization the planned theme for the initial (*i-1*) messages changed with the mobile device, once the latter sends the *i*th message. The protocol instrument Scyther validates the aliveness claim

II. Conclusion

During this paper, we tend to propose a unique authentication theme for mobile cloud computing, Data Digest-based Authentication consists of 3 phases: registration, authentication, and update. With these phases, Data Digest Authentication utilizes hashing, additionally to traditional user id and secret primarily based authentication, to make sure confidentiality and integrity throughout the authentication method. It can survive a range of various attacks, like man-in-the-middle, replay attacks, etc.

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