

## Performance Measurement of WLAN Based On Medium Access Control for Wirelessly Connected Stations

Md. Abbas Ali Khan<sup>1</sup>, Atif Ul Aftab<sup>2</sup>, Md. Imran Hossain Imu<sup>3</sup>

<sup>1</sup>(School of Technology and Health, KTH Royal Institute of Technology, Stockholm, Sweden)

<sup>2</sup>(Lecturer, Department of CSE, Daffodil Institute of IT, Bangladesh)

<sup>3</sup>(Lecturer, Department of CSE, Daffodil Institute of IT, Bangladesh)

**Abstract:** This paper is mainly focuses on the Medium Access Control (MAC) sublayer of the IEEE 802.11 standard for Wireless Local Area Network (WLAN) and delay measurement among the network and also compare of the traffic load. The MAC sublayer is standard mainly includes two fundamental access methods. One is called DCF (Distributed Coordination Function), which uses Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA) approach. Other is PCF (Point Coordination Function), which is based on polling to determine the station that can transmit next.

**Keyword:** DCF, Delay, MAC and PCF.

### I. Introduction:

A wireless LAN is one that makes use of a wireless transmission medium, ad hoc, networking, and coverage of locations difficult to wire. The MAC sub layer in the standard mainly includes two fundamental access methods. One is called DCF (Distributed Coordination Function) which uses Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA) approach. Other is PCF (Point Coordination Function) which is based on the polling to determine the station that can transmit next and provide contention-free service [1]. The physical CSMA/CA, DCF and PCF utilize virtual carrier sense mechanism to determine the state of the medium. This is implemented with the help of Network Allocation Vector (NAV) which provides each station with a prediction of future traffic on the medium. The standard also allows for fragmentation of the MAC data units into smaller frames, which is very favorable in case of wireless channel as it is not reliable enough to transmit longer frames. In this Paper a simulation is shown of a WLAN network with nine wireless stations named as node\_0 up to node\_8. Node\_0 has its **Access Point Functionality** enabled and acting here as Point Coordinator. These stations are configured in four ways as (i) **DCF Scenario:** DCF is implemented in all stations in the wireless local area networks by default. With DCF, every station senses the medium before transmission, defer as long as the medium is busy and when medium goes idle, wait for a random back off interval. After back off interval, if medium is still idle, the station initiates data transmission or optionally exchange RTS (Request to send)/CTS (Clear to send) frames [1]. (ii) **DCF\_PCF Scenario:** PCF access method is implemented 'optionally' in an infrastructure network along with the DCF. With PCF, the access point (AP) in the network acts as a Point Coordinator (PC). The PC uses Polling based approached to determine which station can transmit data. It is optional for the stations to participate in the PCF. Those stations which respond to the poll received from the PC are called CF-Pollable stations. The PCF functionality has been enabled on node\_0, node\_1, node\_3, node\_5 and node\_7. (iii) **DCF\_Frag Scenario:** This scenario has been duplicated from DCF scenario and here on all stations; Fragmentation Threshold size of 256 bytes has been set. So fragmentation of the frames will occur if their size exceeds the fragmentation threshold (which is of 256 bytes). Here in case of DCF and during contention period, all fragments of a single frame will be sent as burst with a single invocation of the DCF medium access procedure.

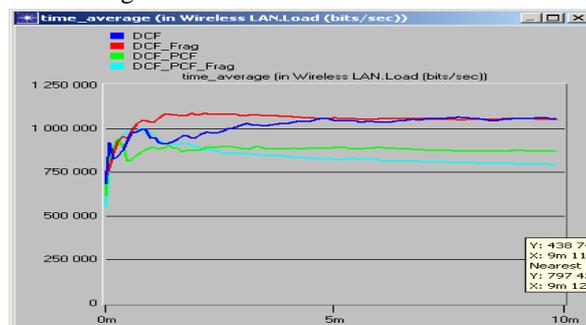


Figure 1.1 Load statistics for DCF and PCF Scenario

(iv) **DCF\_PCF\_Frag Scenario:** This scenario has been duplicated from DCF\_Frag scenario due to which Fragmentation Threshold of 256 bytes has already been set on all stations. Then PCF functionality has been

enabled on node\_0, node\_1, node\_3, node\_5 and node\_7. Here in case of PCF and during a contention free periods, fragments are sent individually following the rules of the PC. As discussed above, DCF is based on CSMA/CA and there is a contention for the medium among the stations. A station only senses the medium before transmission. If the medium is idle, it transmits otherwise if the medium is busy; a station that wishes to transmit has to wait until the current transmission is over and for an additional random back off interval. Thus in this way, there may have collisions which requires retransmission of data after receiving the negative ACKs from receiver [1, 2]. This additional traffic due to the retransmission causes the traffic load in the network. In case of PCF, the existence of AP and CF-Pollable stations creates the contention free period and each station transmits its required data only on its designated interval according to the poll issued by the AP. This causes the collision free transmission and there will be less requirement of retransmission. In the figure 1.1 shown, it has been observed clearly that DCF scenario without PCF (shown in Red and Blue colors) is suffering with higher loads as compared to DCF scenario with PCF (shown in Green and Sky blue colors). The possible cause of this high load is additional traffic of retransmissions, required due to possible collisions in case of DCF scenario. Let see the effect of utilizing PCF and Fragmentation on delay first.

## **II. Effect of PCF on Delay:**

As discussed above, DCF is based on CSMA/CA in which station has to wait until the medium becomes idle and then additional wait for back off time before it can transmit anything. In addition to this, as PCF pollable traffic has more priority so it has to wait until the traffic of any CF-Pollable station becomes over. In addition to this, in DCF, station has to wait for additional SIFS (Shortest IFS) before sending of data and another SIFS before getting an ACK from the receiver [1]. While in case of PCF, station suffers on delay of its polling turn which depends on the number of Pollable stations. If number of pollable stations is higher, then delay for receiving a poll from the PC will be higher. But as here number of pollable stations is only five due to which polling delay is significantly very low.

## **III. Effect of Fragmentation on Delay:**

In fragmentation, the frames which are larger than 256 bytes are fragmented into smaller frames. This will cause additional delay because at the transmission side, first large frame will be chopped down into smaller subframes, then each subframe will be added with its header. Similarly, at the receiver end, header will be removed from each subframe and then reassembled into a larger frame. In case of DCF, this delay can be seen significant because here each packet will have to wait before it can access the medium. Similarly, as the number of frames has been increased, the number of collisions may be increased too. This will cause more retransmissions resulting in more delays for the reassembly of the error free subframes into larger frame at the receiver side. The delay due to this fragmentation is relatively small in case of PCF enabled scenario. The effect of PCF and Fragmentation on Delay can be seen in the figure 1.2. As discussed above, DCF\_Frag (shown with red line) has the highest delay followed by DCF (shown with blue line). DCF\_PCF\_Frag (shown with sky blue color) has fragmentation delay and the DCF\_PCF (shown with Green color) has the smallest delay as compared to other scenarios.

## **IV. Effect of PCF on Throughput:**

The scenario without PCF will suffer more delays and thus throughput is small as compared to the scenario with PCF. Secondly, because of the retransmissions due to any possible collisions, effective throughput will also be relatively small as compared to scenario with PCF. Effect of Fragmentation on Throughput, refer to the discussion above, in case of DCF only, the number of frames will be increased, which will cause more collisions resulting in greater number of retransmissions. This will cause effective throughput to be decreased as compared to the throughput in case of scenario with PCF. The above two findings can be seen clearly in the figure 1.3. The highest throughput is in DCF\_PCF (shown with Green color) followed by DCF\_PCF\_Frag (shown with sky blue color). The low throughput is in DCF (shown with Blue color) followed by the DCF\_Frag (shown with red color) showing the lowest throughput among the four scenarios. Referred to the figure 1.4 and figure 1.5, in DCF\_PCF scenario, node\_2 has its PCF Functionality disabled while node\_3 has its PCF functionality enabled. Figure 1.4 shows the delay in node\_2 and node\_3 while figure 1.5 shows the retransmission attempts in node\_2 and node\_3. As seen in figure 1.4, node\_2 and

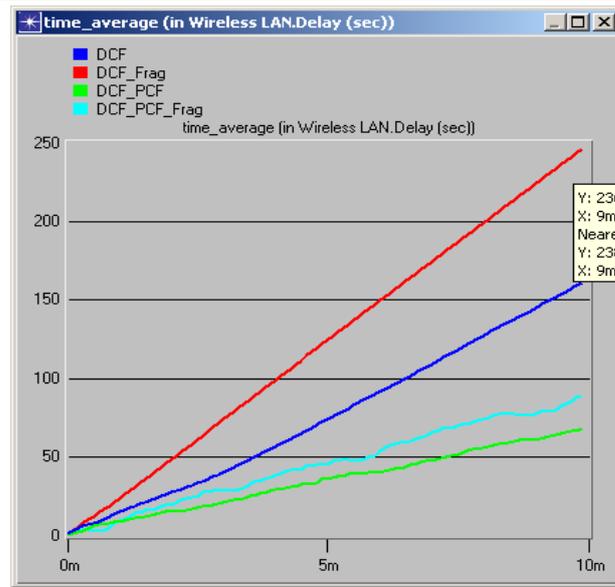


Figure 1.2 Delay statistics for all four scenarios

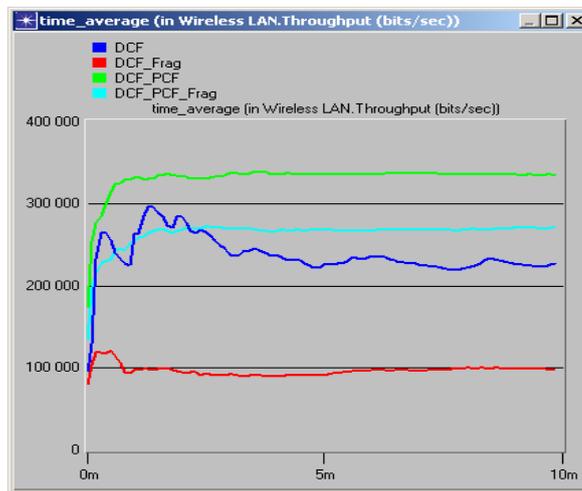
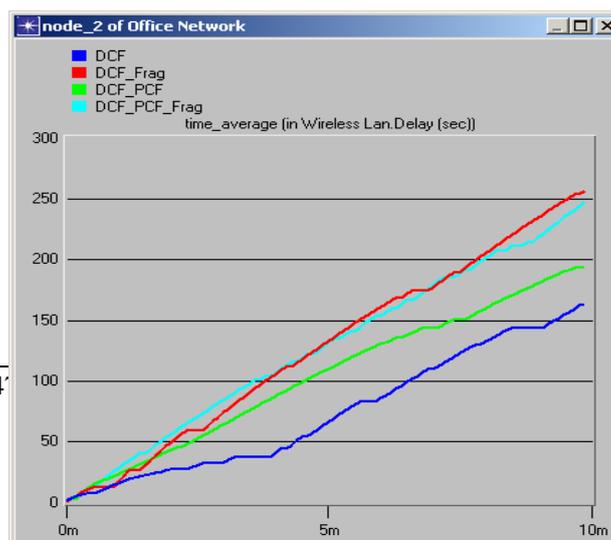


Figure 1.3 Throughput statistics for all four scenarios

node\_3 has almost the same results for DCF and DCF\_Frag scenarios. But in case of DCF\_PCF and DCF\_PCF\_Frag, the delay is much higher in node\_2 as compared to that in node\_3. The reason for this is that in case of PCF enabled in some nodes, the CF-pollable nodes (like node\_3) can transmit right away without any delay in the collision free period (CFP) and the non-CF-Pollable nodes (like node\_2) experience higher delays due to the existence of the collision-free period during which only CF-Pollable nodes are permitted to transmit. In the figure 1.5, the number of retransmissions in node\_3 is considerably low as compared to that in node\_2 for the four scenarios. The reason for this is that, in node 3, when PCF is used, it guarantees that there will be no collision as each station is transmitting at



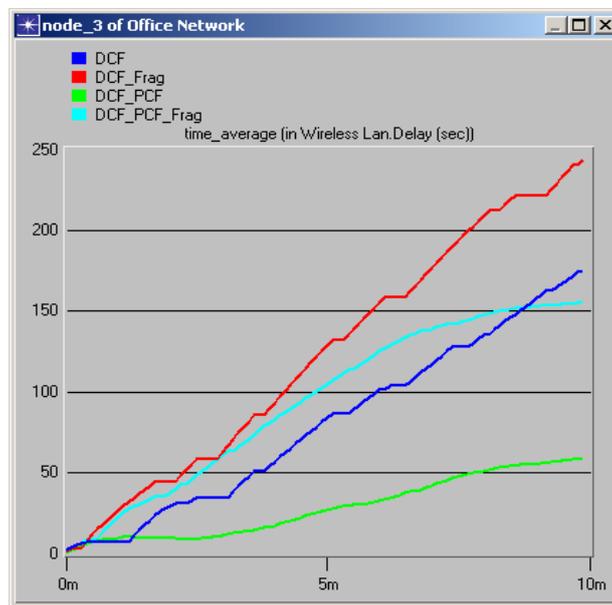
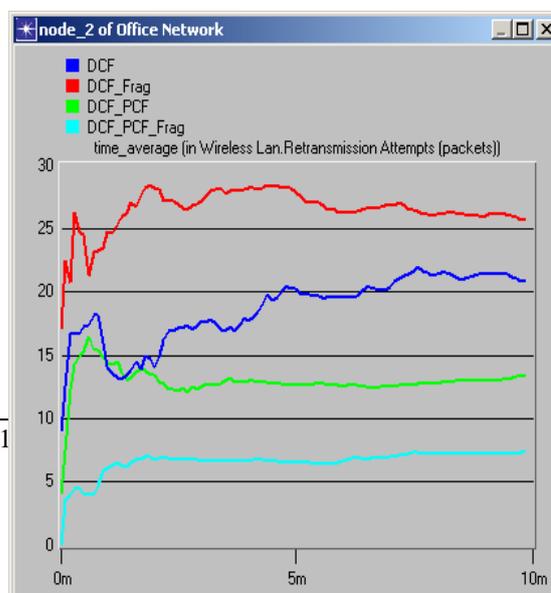


Figure 1.4 Delay Station for node\_2 and node\_3 (in DCF\_PCF scenario)

its desired interval and thus there will be no retransmission. So the number of transmissions will be lower for DCF\_PCF (shown with green color) and DCF\_PCF\_Frag (shown with sky blue color) in node\_3 as compared to node\_2. Secondly, the nodes with PCF enabled do not compete for access to the medium, so there will be lower number of stations (nodes) that will contend and transmit during contention period. Thus the number of retransmissions for DCF\_Frag (shown with red color) and DCF (shown with blue color) will also be relatively smaller in node\_3 as compared to those in node\_2. Referred to the figure 1.6 and 1.7 are actually doing comparison between the two scenarios where in one scenario, have only two PCF enabled nodes ( called DCF\_two PCF) and PCF enabled on all Referred to the Figure delay statistics for

the other scenario when have nodes (called DCF\_allPCF). 1.6 is shown the comparison of DCF\_twoPCF and DCF\_allPCF.



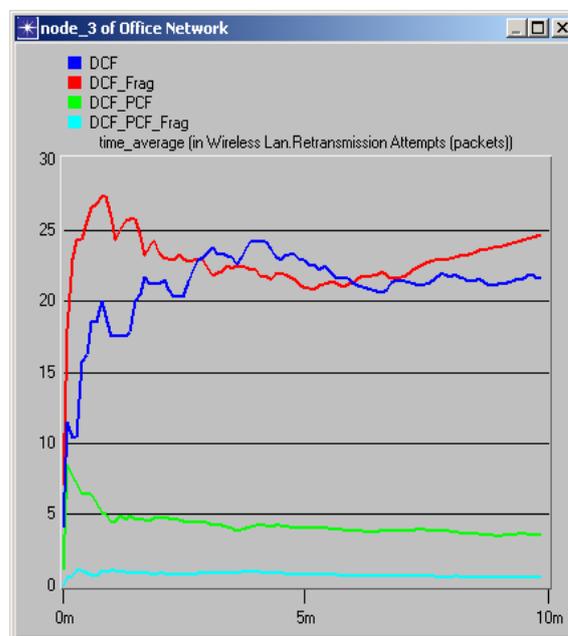


Figure 1.5 Retransmission Attempts statistics for node2 and node 3 (in DCF\_PCF scenario)

It can be seen that the delay in DCF\_twoPCF (shown with Pink color) is much higher as compared to the delay in case of DCF\_allPCF (shown with Yellow color). As discussed above, the reason for it is that have much chance of collisions and the respective retransmissions when have only two PCF enabled nodes, causing much delay. But when we have all PCF enabled nodes, then there are almost no collisions and retransmissions resulting in much less delay. For the same reason, in Figure 1.6 (a) throughput in DCF\_allPCF (shown with Yellow color) is much higher as compared to that DCE\_twoPCF (pink color).

According to the Figure 1.7 shown, is getting unusual result of load comparison between DCF\_allPCF and DCF\_twoPCF, because usually if it has less delay and more throughputs, so expect fewer loads in the network. But here DCF\_allPCF have more loads (shown with Yellow color) as compared to that in case of DCF\_twoPCF (shown with pink color). The possible reason for it can be that, in case of DCF\_allPCF, the number of Polls sent by the PC is higher due to the greater number of CF-Pollable nodes as compared to the number of polls sent when only two CF-Pollable nodes have in DCF\_twoPCF scenario. This can cause extra load in the network even if there are not any collisions and the resulting retransmissions. If the number of CF-Pollable

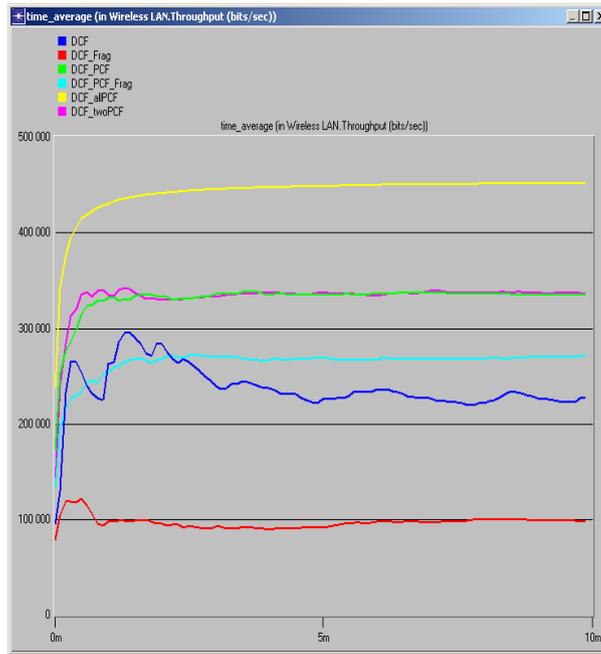


Figure 1.6(a)

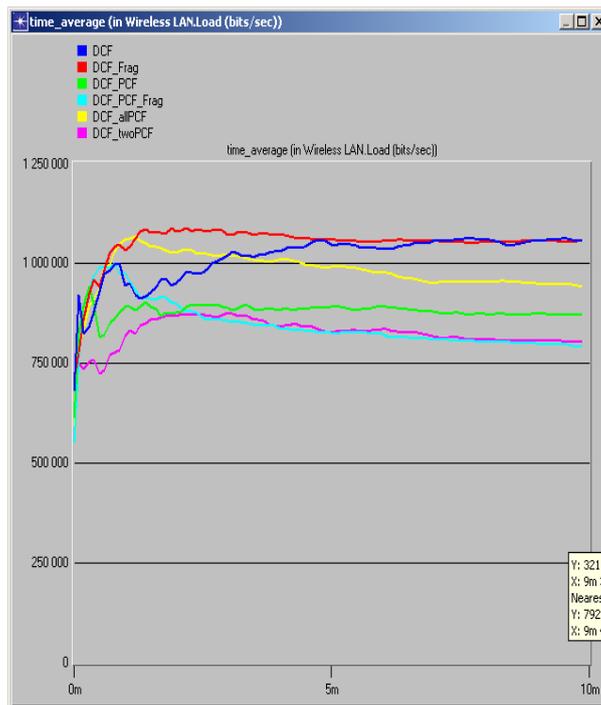


Figure 1.6(b)

Figure 1.6 (a) Throughput and 1.6(b) Delay statistics for DCF\_compared to that in DCF\_twoPCF (shown with Pink color),allPCF and DCF\_twoPCF

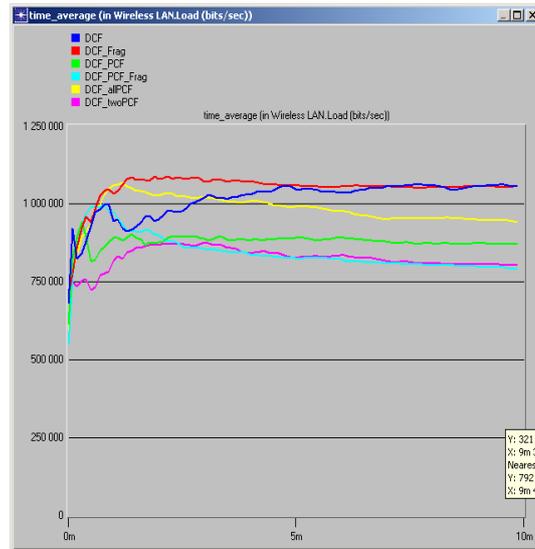


Figure 1.7 Load statistics for DCF\_allPCF and DCF\_twoPCF

nodes in the network is increased, this load effect can be experienced more significantly in that case. The Medium Access Delay for all six scenarios is shown in Figures 1.8 and 1.9. From the Figure 1.8, it can be seen clearly that DCF\_Frag (shown with Green color) has the highest medium access delay followed by the DCF (shown with Red color) as compared to remaining PCF enabled scenarios. As discussed in above, the reason for high delay in DCF scenarios (DCF\_Frag and DCF) is that there is no PCF enabled nodes and all the nodes are accessing the medium during the contention period resulting in collisions and corresponding retransmissions. Between DCF\_Frag and DCF, DCF\_Frag has more delay because of the fragmentation issues as discussed above. For the remaining four PCF enabled scenarios, results are seen unexpected in Figure 1.8 so their simulation results have been generated again and shown in Figure 1.9. As per Figure 1.9 shown, it can be seen clearly that DCF\_PCF\_Frag (shown with sky blue color) has higher medium access delay as compared to the remaining DCF\_twoPCF, DCF\_PCF and DCF\_allPCF scenarios. The reason for this is the same fragmentation issues as discussed in above. For the remaining three scenarios, DCF\_twoPCF (shown with pink color) has more delay compared to DCF\_PCF and DCF\_allPCF. The reason for it is the effect on delay due the number of CF-Pollable nodes in the network discussed in above. In DCF\_twoPCF, there are only two CF-Pollable nodes resulting in higher number of nodes to work in Contention period resulting in more collisions and corresponding retransmissions. In DCF\_PCF (shown with green color), almost half of all the nodes are CF-Pollable resulting in half of the nodes in the network to work during Contention Free period resulting in significantly less number of collisions and retransmissions. Finally in DCF\_allPCF (shown with Yellow color), the medium access delay is the lowest as almost all the nodes are working in Contention Free period and there are not any significant collisions and retransmissions.

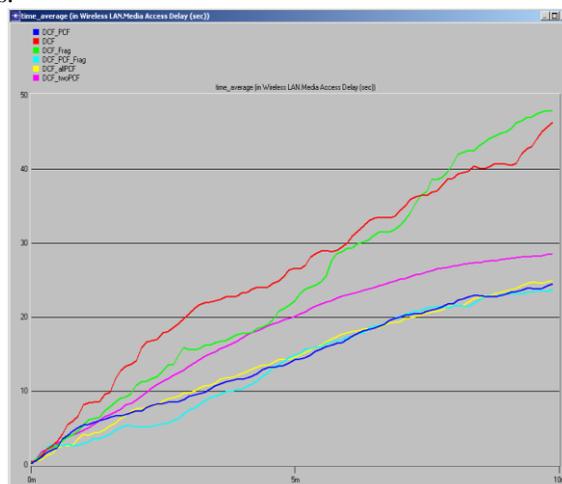


Figure 1.8 Medium Access Delay statistics for DCF\_Frag and DCF

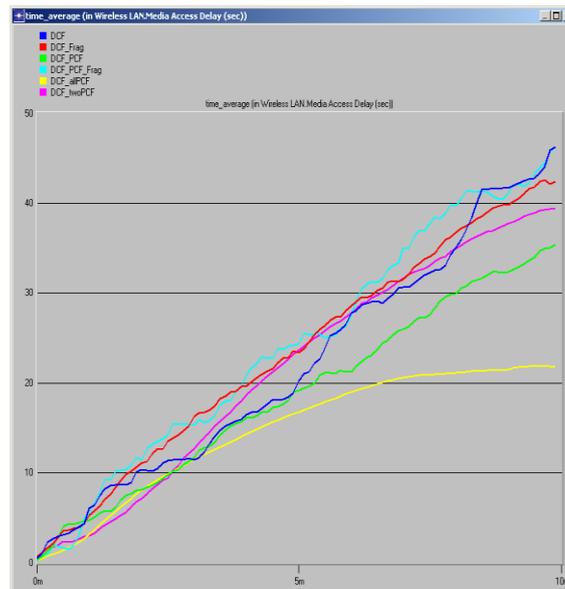


Figure 1.9 Medium Access Delay statistics for DCF\_PCF\_Frag, DCF\_twoPCF, DCF\_PCF and DCF\_allPCF

## V. Conclusion

The DCF\_Frag has the highest delay followed by DCF when PCF is not enabled on all the nodes in the network. When PCF is enabled in the network, DCF\_PCF\_Frag has higher delay followed by DCF\_twoPCF, DCF\_PCF and DCF\_allPCF respectively. Polling delay is happened if the number of pollable station is more. To make less such delay uses less number of polling stations. And traffic load occurs due to retransmission. Less number of collisions and retransmissions are occurred during the contention free period. If it is being ensured less traffic load (retransmission) in the network then potentiality is more to get less delay and more throughputs.

## VI. References:

- [1] Stallings William, Wireless Communication & Networks. 2<sup>nd</sup> edition, P-435, prentice- Hall of India, New Delhi-110 0001
- [2] Larry L. Peterson & Bruce S. Davie, Computer Networks: A System Approach, 4<sup>th</sup> edition, section 2.8.2, P-140