

“Software Theft Detection for JavaScript programs based on dynamic birthmark extracted from runtime heap graph”

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Abstract: Software's, programs are valuable assets to developer companies. However, the source code of programs can be theft and JavaScript programs whose code is easily available which is a serious threat to the industry. There are techniques like watermarking which prove the ownership of the program but it can be defaced and encryption which changes source code but it may decrypted and also it cannot avoid the source code being copied. In this paper, we use a new technique, software birthmark, to help detect code theft of software or program. A birthmark is a unique characteristic of a program which is used to identify the program. We extract the birthmark of software from the run-time heap by using frequent sub graph mining and search the same in suspected program.

Keywords: Heap graph, software birthmark, frequent sub graph mining.

I. Introduction

Software theft is a serious issue in the industry. Software theft detection is very important for software industry. The watermarking is a solution used to prove ownership. Another approach is encryption in which a source code is transformed in such way that it becomes more difficult to understand. A new technique for software theft detection is software birthmark which does not insert any code to the software and it depends on the characteristics of the programs. There are two types of software birthmarks, static birthmarks and dynamic birthmarks. Static birthmarks are extracted from the syntactic structure of a program. Dynamic birthmarks are extracted from the dynamic behavior of a program at run-time. We use dynamic birthmark approach for software theft detection.

A birthmark is used to identify software theft, to detect software theft; the dynamic birthmark of the original program is first extracted by using frequent sub graph mining. The same software birthmark will be search in suspected program. If the birthmark is found, the suspected program is a copy of the original program.

II. Related work

- [1] A. Monden, H. Iida, K. I. Matsumoto, K. Inoue, and K. Torii have proposed Watermarking based software theft detection method for java programs.
- [2] C. Collberg, E. Carter, S. Debray, A. Huntwork, J. Kececioglu, C. Linn, and M. Step have proposed Dynamic path-based software watermarking for software theft detection.
- [3] X. Wang, Y.-C. Jhi, S. Zhu, and P. Liu have proposed dynamic behavior based software theft detection
- [4] G. Myles and C. Collberg have detected software theft via whole program path birthmarks.
- [5] D. Schuler, V. Dallmeier, and C. Lindig have shown robust dynamic birthmark for java program software theft detection.
- [6] H. Tamada, K. Okamoto, M. Nakamura, A. Monden, and K. I. Matsumoto have designed and evaluated dynamic software birthmarks based on API Calls.
- [7] P. Chan, L. Hui, and S. Yiu have proposed dynamic JavaScript birthmark based on the run-time heap (JsBirth).
- [8] P. P. F. Chan, L. C. K. Hui, and S.M. Yiu have used heap memory analysis for dynamic software birthmark formation and comparison for java programs.

III. Proposed algorithm

Based on above discussion proposed work is to implement and analyze birthmark based software theft detecting system for JavaScript programs. Specifically the proposed work is stated below.

1. **Design and implementation of the graph generator and filter-**The JavaScript heap profiler will run a JavaScript program and takes multiple heap snapshots in the course of its execution. The graph generator and filter will traverse the objects in the heap snapshots and builds heap graphs out of them. Heap Graph generates for each heap snapshot. Heap snapshot consist arrangement of objects which are created runtime. Heap graph contains various types of objects and edges. Graph filter is used to filter unwanted objects and

edges. Object of type OBJECT and CLOSURE are considered and remaining types of objects are discarded. References of type ELEMENT and PROPERTY are only included other are discarded.

2. **Design and implementation of the graph merger** –The graph merger will merge the filtered heap graphs together to form one single graph. According to object id all heap graphs are combined together. We merge all the graphs one by one by taking the union set of the nodes and edges of the two graphs being merged. In order to make the resulting superimposition graph also connected, it need to ensure that there is at least one object in common (with the same object ID) in two graphs before superimposing them.
3. **Design and implementation of modified Graph Selector-** The sub graph selector will select a sub graph from the heap graph to form the birthmark of the original program. The modified Graph Selector will use frequent sub graph mining to get the frequent sub graph that appears in all the heap graphs that are extracted from the program by graph generator and filter. The proposed modified Graph Selector will be tested for its suitability in using it as the more representative birthmark of the program.
4. **Design and implementation of the detector** –The detector will search for the birthmark of the original program in the heap graph of the suspected program.
5. **Application and analysis** of above mentioned scheme for its usefulness, robustness and accuracy in protecting the intellectual property rights of JavaScript developers.

IV. Simulation Results

Module 1:

1) Graph generator and filter:

Algorithm:

Input: JavaScript heap profile from JavaScript heap profiler

Output: A set of filtered heap graphs captured at different points of time with the annotated nodes in the heap by its object ID.

I. Graph generator:

For each snapshot taken using the Chromium browser, perform a depth first search traversal of it and print out the heap graph with annotated nodes by its object ID (objects) and edges (reference) and then pass it to a filter.

II. Graph filter:

Graph Filter working is as follows.

For each JavaScript Objects we do following:

- I. If Object is of type **INTERNAL** or **ARRAY** or **STRING** or **CODE** discard the object.
- II. If Object is of type **OBJECT** or **CLOSURE** include only **ELEMENT** and **PROPERTY** references and
- III. Filter out every other reference.

Module 2:

1) Graph merger:

Algorithm:

Input: Multiple labelled connected heap graphs G from Graph generator and filter

Output: A superimposition of multiple labelled connected heap graphs G into a one single graph M that includes all the nodes and edges appearing in the input heap graphs.

The above algorithm superimposes all the graphs one by one by taking the union set of the nodes and edges of the two graphs being merged. In order to make the resulting superimposition graph also connected, it need to ensure that there is at least one object in common (with the same object ID) in two graphs before superimposing them.

2) Sub graph selector:

Algorithm:

Input: A set of filtered heap graphs captured at different points of time with the annotated nodes in the heap by its object ID (obtained from module 2- graph generator and filter).

Output: the frequent sub graphs to be used as the birthmark

The modified Graph Selector uses frequent sub graph mining to get the frequent sub graph that appears in all the heap graphs that are extracted from the program by graph generator and filter. For frequent sub graph mining gSpan algorithm is used which is complete frequent sub graph mining algorithm on labeled graphs. gSpan builds a new lexicographic order among graphs, and maps each graph to a unique minimum DFS code as its canonical label. Based on this lexicographic order, gSpan adopts the depth first search strategy to mine frequent connected sub graphs efficiently.

This module uses java library ParSeMiS (Parallel and Sequential Graph Mining Suite). The library searches for frequent, interesting substructures in graph databases. It has implementation of gSpan algorithm for frequent sub graph mining.

3) Sub graph detector:

1. Take birthmarks obtained from plaintiff program. Take single large merged graph obtained from suspected program.
2. For each birthmark
 - a. Check if birthmark is present in single large merged graph.
 - b. If birthmark is found report “yes”.
3. If no birthmark is found report “not found”.

Algorithm:

Input: sub graphs from the original program and the entire heap graph of the suspected program

Output: Determines whether the selected sub graphs of the plaintiff program can be found in the heap graph of the suspected program.

- I. Takes sub graphs of the objects under the Window objects from the suspected program
- II. Use sub graph monomorphism to check whether the subgraphs of the plaintiff program can be found in them.
- III. If a match is found, raise an alert and reports where the match is found.

Experimental Result

We proposed five modules in this paper, three modules are implemented and result is described in following section. Sub graph detector and analysis are not implemented. The results of first three modules are mentioned

We took first heap snapshot of www.google.com and passed it to the Graph Generator. It creates Heap graph by considering all types of object and all types of edges the result of first heap snapshot is as follows

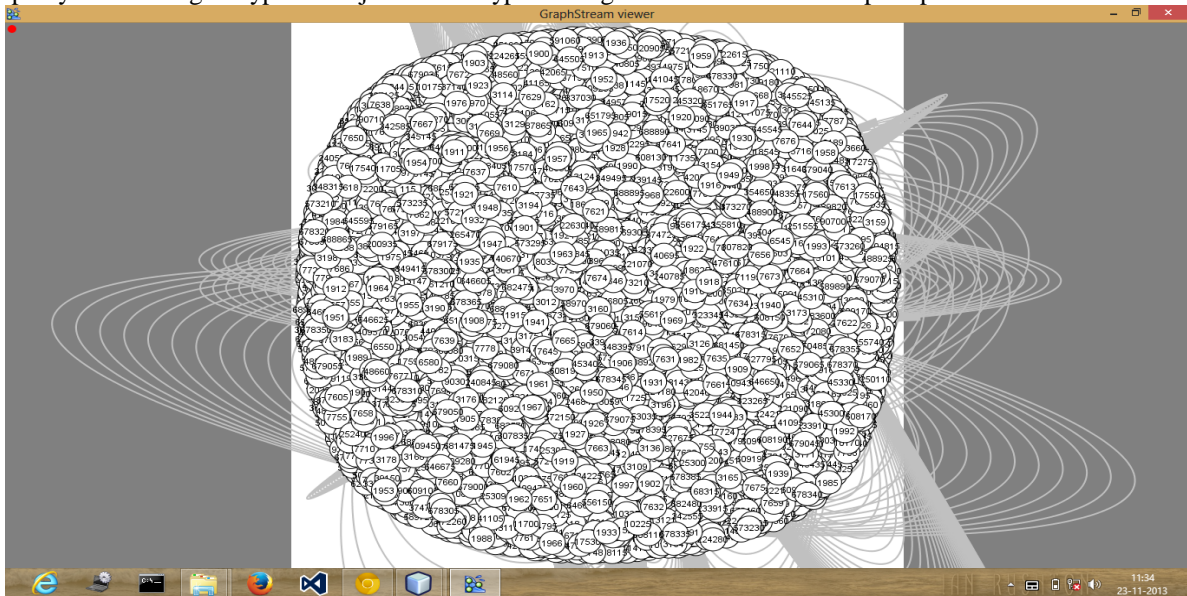


Figure 1: Results of Graph Generated by first module from sample firstheapdump

We took second heap snapshot of www.google.com and passed it to the Graph Generator. It creates Heap graph by considering all types of object and all types of edges the result of second heap snapshot is as follows

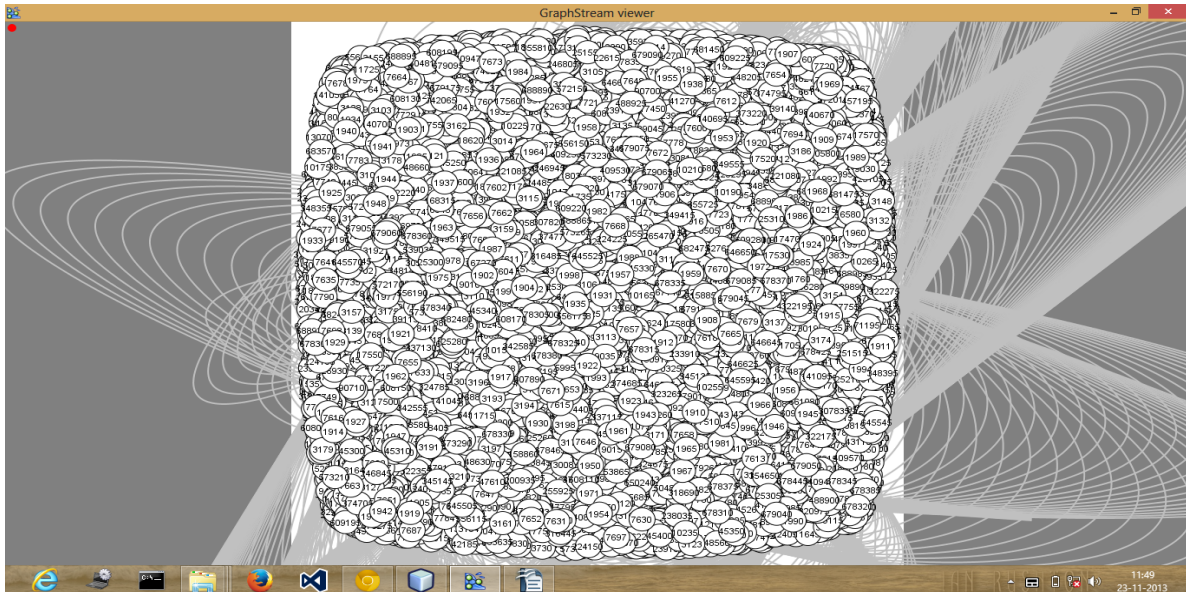


Figure 2: Results of Graph Generated by first module from sample second heapdump

First heap snapshot of www.google.com passed to the Graph Generator which generates heap graph as shown in figure 1. Then resulting graph passed as input to the Graph Filter which only consider Object of type OBJECT and CLOSURE and remaining types of objects are discarded. References of type ELEMENT and PROPERTY are only included other are discarded. The resultant filtered graph is as follows

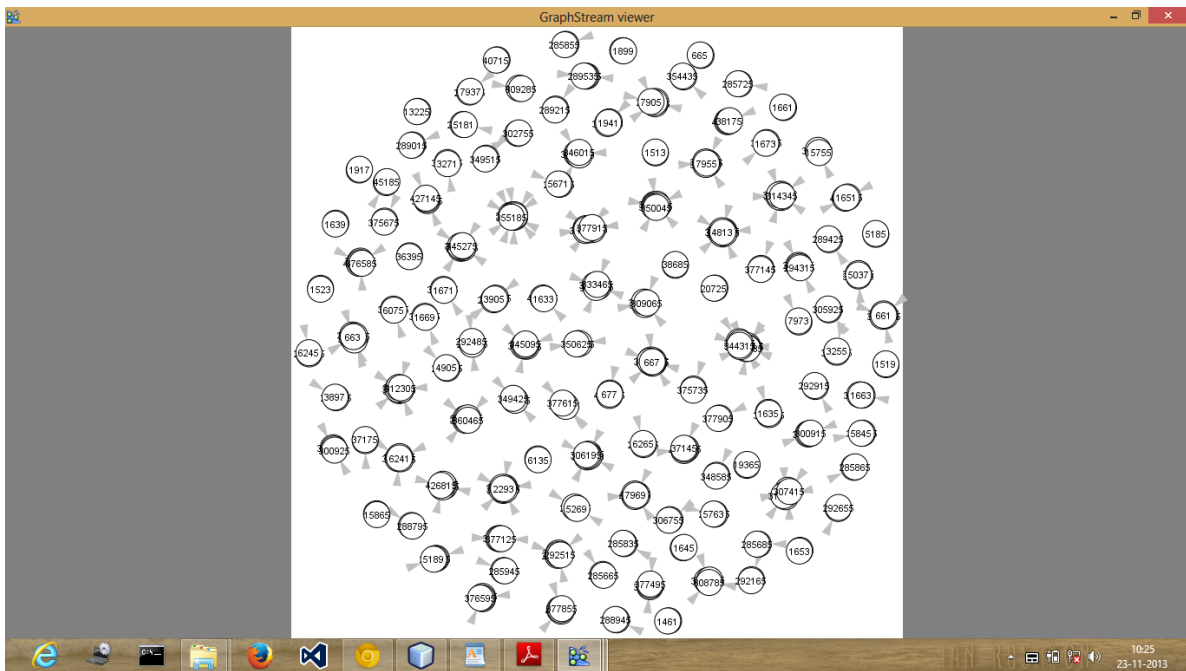


Figure 3: Results of Graph filtered of graph from fig.1 by second module

Second heap snapshot of www.google.com passed to the Graph Generator which generates heap graph as shown in figure 2. Then resulting graph passed as input to the Graph Filter which only consider Object of type OBJECT and CLOSURE and remaining types of objects are discarded. References of type ELEMENT and PROPERTY are only included other are discarded. The resultant filtered graph is as follows

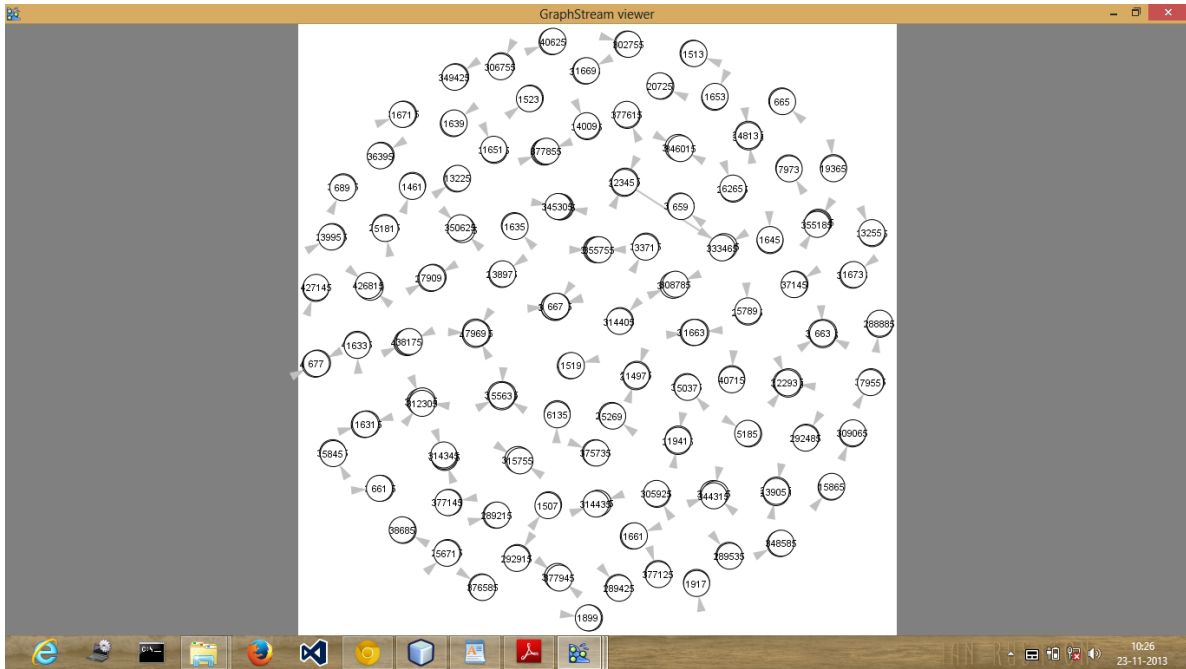


Figure 4: Results of Graph filtered of graph from fig.2 by second module.

In the heap graph all objects are assigned a unique id. Two snapshots which are shown in figure 3 and figure 4 are passed to the Graph Merger which produces single graph and merging done on the basis of object id. The resultant single graph is as follows

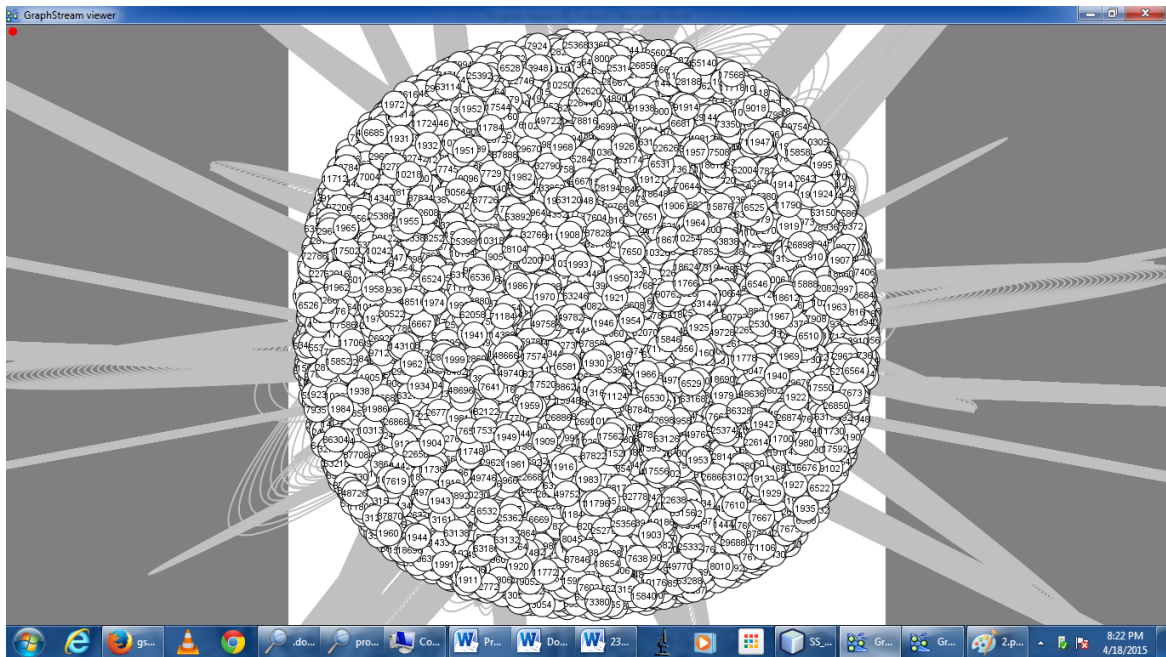


Figure 5: Results of merging of two Graphs from fig.3 and fig 4

Frequent sub graph Mining is applied on single graph generated by Graph Merger as shown in figure 5 which searches subgraph with maximum occurrence in the largest graph. The resultant subgraph of graph displayed in figure 5 is as follows

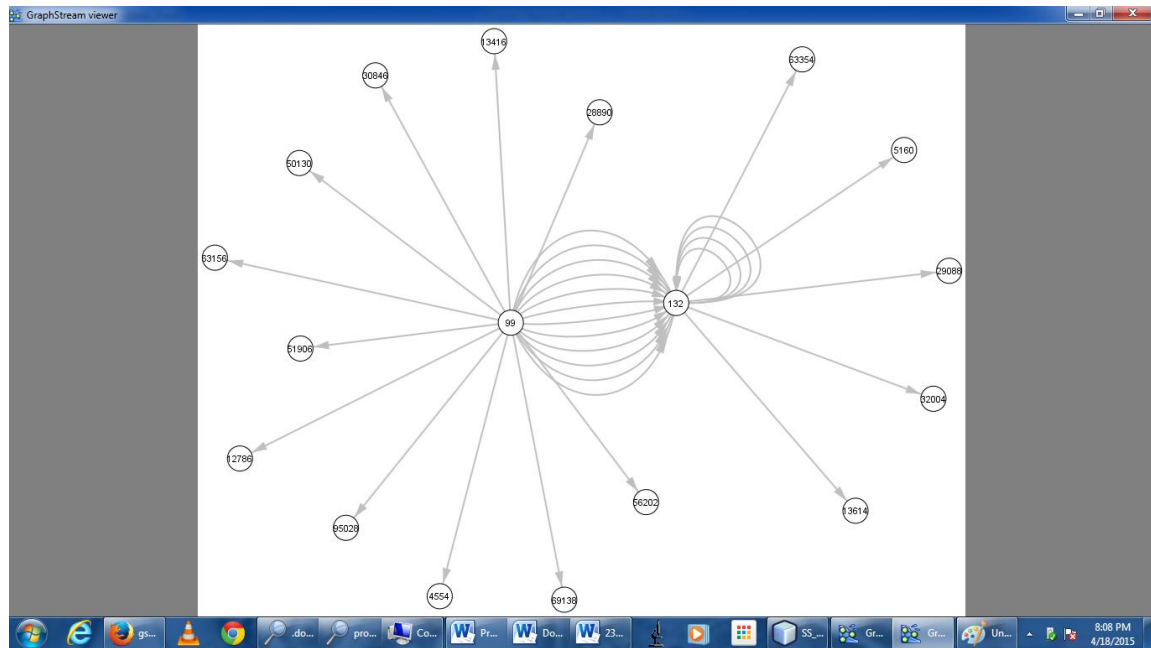


Figure 6: Sub graph selected by sub graph selector module as a software birthmark

The resultant sub graph of graph displayed in figure 6 will be search in suspected program. If birthmark will found it will display message as birthmark found otherwise it will display as follows in fig.7

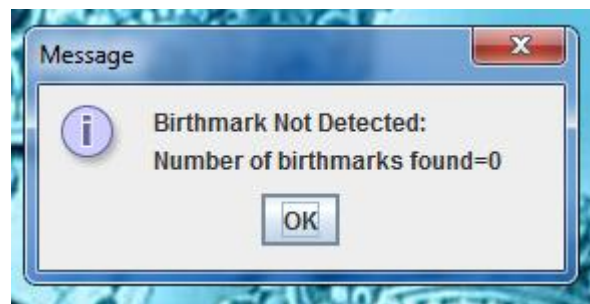


Fig 7: Result display after birthmark searching

V. Conclusion

We proposed a robust heap graph based software birthmark system for JavaScript programs. We evaluated our birthmark system using 50 large-scale websites and the experiment results are promising with 100% accuracy

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