Discovering Frequent Patterns with New Mining Procedure

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Abstract: Efficient algorithm to discover frequent pattern are crucial in data mining research. Finding frequent itemsets is computationally the most expensive step in association rule discovery .To address these issues we discuss popular techniques for finding frequent itemsets in efficient way. In this paper we provide the survey list of existing frequent itemsets mining techniques and proposing new procedure which having some advantages by comparing with the other algorithms.

Keywords—Association rules, data mining, frequent itemsets, FPM, minimum support.

I. INTRODUCTION

Data mining also known as Knowledge Discovery and Database (KDD) has been perceived as a new research is for database research. The aim of data mining is to automate the process of finding interesting patterns and trends from a given data. Discovering of association rules is an important data mining task. A procedure so called mining frequent itemsets is a core step of association rules discovering introduce in [1].

Past transaction analysis is a commonly used approach in order to improve the quality of such decision. As pointed in [2], the progress in bar-code technology has made it possible for retail organizations to collect the store massive amount of sales data so called Basket data that stores item purchased per transaction basis. Most existing algorithm, therefore, are focused on mining frequent itemsets [7]. Several algorithm have been proposed to mine frequent itemsets. A classical algorithm is Apriori introduced in[1]. Variation of Apriori were proposed so far such as hash-based algorithm [3], a dynamic itemset counting technique (DIC) [4] and a Partition algorithm. In addition, many research have been developed algorithms using tree structure, such as H-mine and FP-growth [5].

There are two main strategies for mining frequent itemsets: the candidate generation-and-test approach and the pattern growth approach. Apriori [2] and its several variation belong to the first approach, while FP-growth [5] and H-mine are example of the second. Apriori algorithm [2] suffer from

the problem spending much of their time discard the infrequent candidate on each level. The FP-growth (Frequent Pattern -Growth) [5] algorithm differs basically from the level-wise algorithm, that use a "candidate generate and test" approach.

II. PROBLEM STATEMENT

Formally as defined in [1], the problem of mining association rules is started as follows: Let I = $\{i_{1}, i_{2}, \dots, i_{M}\}$. Let D be a set of transactions, where each transaction T is a set of item such that $T \subseteq I$. Associated with each transaction is unique identifier, called its transaction id is TID. A transaction T contains X, a set of some items in I, if $X \subseteq T$ An association rule is amplification of the form $X \Rightarrow Y$ where $X \subset I$, $Y \subset I$ and $X \cap Y = \emptyset$. The meaning of such a rule is that transactions in database, which contain the items in X, tend to also contain the items in Y. The rule $X \subseteq Y$ holds in the transaction set D with confidence c if among those transactions that contains X c% of them also contain Y. The rule $X \Rightarrow Y$ has support s in the transaction set D if s% of transactions in D contain X UY. The problem of mining association rules that have support and confidence greater than the user-specified minimum support and minimum confidence respectively. Conventionally, the problem of discovering all association rules is composed of the following two steps: 1) Find the large itemsets that have transaction support above a minimum support and 2) From the discovered large itemsets generate the desired association rules. The overall performance of mining association rules can be achieved by using the first step. After the identification of large itemsets, the corresponding association rules can be derived in a straightforward manner. In this paper, we have developed a method to discover large itemsets from the transaction database, thus finding a solution for the first sub problem.

A. Algorithms For Mining Frequent Itemsets

In this each row of database represents a transaction which has a transaction identifier (TID), followed by a set of items. An example of horizontal layout dataset is shown in Table 1. In vertical layout data set, each column corresponds to an item, followed by a TID list, which is the list of rows that the item appears. An example of vertical layout database set is as shown in diagram for the Table 2.

TID	ITEMS
T1	11, 12, 13, 14, 15, 16
T2	I1, I2, I4, I7
T3	I1, I2, I4, I5, I6
T4	I1, I2, I3
T5	13, 15

 Table 1: Horizontal layout based database

Table 2 : Vertical layout based database

ITEM	TID_list
I1	T1, T2, T3, T4
I2	T1, T2, T3, T4
I3	T1, T4, T5
I4	T1, T2, T3
15	T1, T3, T5
I6	T1, T3
I7	T2

(1). Apriori Algorithm

The first algorithm for mining all frequent itemsets and strong association rules was the AIS algorithm. Shortly after that, the algorithm was improved and renamed Apriori. Apriori algorithm is, the most classical and important algorithm for mining frequent itemsets. Apriori is used to find all frequent itemsets in a given database DB. The key idea of Apriori algorithm is to make multiple passes over the database. It employs an iterative approach known as a breadth-first search (level-wise search) through the search space, where k-itemsets are used to explore (k+1)-itemsets.

The working of Apriori algorithm is fairly depends upon the Apriori property which states that" All nonempty subsets of a frequent itemsets must be frequent" [1]. It also described the anti monotonic property which says if the system cannot pass the minimum support test, all its supersets will fail to pass the test [1],[2]. Therefore if the one set is infrequent then all its supersets are also frequent and vice versa. This property is used to prune the infrequent candidate elements. In the beginning, the set of frequent 1-itemsets is found. The set of that contains one item, which satisfy the support threshold, is denoted by L_1 . In each subsequent pass, we begin with a seed set of itemsets found to be large in the previous pass. This seed set is used for generating new potentially large itemsets, called candidate itemsets, and count the actual support for these candidate itemsets during the pass over the data. At the end of the pass, we determine which of the candidate itemsets are actually large (frequent), and they become the seed for the next pass. Therefore, L_1 is used to find L_2 , the set of frequent 2-itemsets, which is used to find L_3 , and so on, until no more frequent k-itemsets can be found. The feature first invented by [1] in Apriori algorithm is used by the many algorithms for frequent pattern generation.

The basic steps to mine the frequent elements are as follows [2]:

- Generate and test: In this first find the 1-itemset frequent elements L₁ by scanning the database and removing all those elements from C which cannot satisfy the minimum support criteria.
- Join step: To attain the next level elements C_K join the previous frequent elements by self join i.e. $L_{K-1} * L_{K-1}$ known as Cartesian product of L_{K-1} . i.e. This step generates new candidate k-itemsets based on joining L_{K-1} with itself which is found in the previous iteration. Let C_K denote candidate k-itemset and L_1 be the frequent k-itemset.
- Prune step: C_K is the superset of L_1 so members of C_K may or may not be frequent but all K-1 frequent itemsets are included in C_K thus prunes the C_K to find K frequent itemsets with the help of Apriori property. i.e. This step eliminates some of the candidate k-itemsets using the Apriori property. A scan of the database to determine the count of each candidate in C_K would result in the determination of L_1 (i.e., all candidates having a count no less than the minimum support count are frequent by definition, and therefore belong to L_1). C_K , however, can be huge, and so this could involve grave computation. To shrink the size of C_K , the Apriori property is used as follows. Any (k-1)-itemsets that is not frequent cannot be a subset of a frequent k-itemsets. Hence, if any (k-1)-subset of candidate k-itemsets is not in L_{K-1} then the candidate cannot be frequent either and so can be removed from C_K . Step 2 and 3 is repeated until no new candidate set is generated.

It is no doubt that Apriori algorithm successfully finds the frequent elements from the database. But as the dimensionality of the database increase with the number of items then:

• More search space is needed and I/O cost will increase.

• Number of database scan is increased thus candidate generation will increase results in increase in computational cost.

Therefore many variations have been takes place in the Apriori algorithm to minimize the above limitations arises due to increase in size of database. These subsequently proposed algorithms adopt similar database scan level by level as in Apriori algorithm, while the methods of candidate generation and pruning, support counting and candidate representation may differ.

The algorithms improve the Apriori algorithms by:

- Reduce passes of transaction database scans.
- Shrink number of candidates.
- Facilitate support counting of candidates.

(2). Direct Hashing And Pruning (DHP) Algorithm

It is absorbed that reducing the candidate items from the database is one of the important task for increasing the efficiency. Thus a DHP technique was proposed [3] to reduce the number of candidates in the early passes C_k for k > 1 and thus the size of database. In this method, support is counted by mapping the items from the candidate list into the buckets which is divided according to support known as Hash table structure. As the new itemset is encountered if item exist earlier then increase the bucket count else insert into new bucket. Thus in the end the bucket whose support count is less the minimum support is removed from the candidate set.

In this way it reduce the generation of candidate sets in the earlier stages but as the level increase the size of bucket also increase thus difficult to manage hash table as well candidate set.

(3). FP-Growth Algorithm

FP-tree algorithm [5] is based upon the recursively divide and conquers strategy; first the set of frequent 1-itemset and their counts is discovered. With start from each frequent pattern, construct the conditional pattern base, then its conditional FP-tree is constructed (which is a prefix tree.). Until the resulting FP-tree is empty, or contains only one single path. (Single path will generate all the combinations of its sub-paths, each of which is a frequent pattern). The items in each transaction are processed in L order. (i.e. items in the set were sorted based on their frequencies in the descending order to form a list).

The detail step is as follows: [5]

FP-Growth Method: Construction of FP-tree

- Create root of the tree as a "null".
- After scanning the database D for finding the 1-itemset then process the each transaction in decreasing order of their frequency.
- A new branch is created for each transaction with the corresponding support.
- If same node is encountered in another transaction, just increment the support count by 1 of the common node.
- Each item points to the occurrence in the tree using the chain of node-link by maintaining the header table.

After above process mining of the FP-tree will be done by Creating Conditional (sub) pattern bases:

- Start from node constructs its conditional pattern base.
- Then, Construct its conditional FP-tree & perform mining on such a tree.
- Join the suffix patterns with a frequent pattern generated from a conditional FP-tree for achieving FP-growth.
- The union of all frequent patterns found by above step gives the required frequent itemset.
- In this way frequent patterns are mined from the database using FP-tree.

Definition: Conditional pattern base [5]

A "sub database" which consists of the set of prefix paths in the FP-tree co-occurring with suffix pattern.eg for an itemset X, the set of prefix paths of X forms the conditional pattern base of X which co-occurs with X.

Definition: Conditional FP-tree [5]

The FP-tree built for the conditional pattern base X is called conditional FP-tree.

(4). Frequent Item Graph (FIG) Algorithm

The FIG [7] algorithm is a novel method to find all the frequent itemsets quickly. It discovers all the frequent itemsets in only one database scan. The construction of the *graphical structure* is done in two phases, where the first phase requires a full I/O scan of the dataset and the second phase requires only a full scan of frequent 2-itemsets. The first initial scan of the database identifies the frequent 2-itemsets with a minimum support. The goal is to generate an ordered list of frequent 2-itemsets that would be used when building the *graphical structure* in the second phase.

The first phase starts by arranging the entire database in the alphabetical order. During the database scan the number of occurrences of frequent 2-itemsets is determined and infrequent 2-itemsets with the support less than the support threshold are weeded out. Then the frequent 2-itemsets are ordered in the alphabetical order. Phase 2 of constructing the *graphical structure* requires a complete scan of the ordered frequent 2-itemsets. The ordered frequent 2-itemsets are used in constructing the *graphical structure*. *Advantages:*

- The quick mining process that does not use candidates.
- Mining the set of all frequent itemsets in the database by scanning the entire database only once.

• I/O costs spent in mining the set of all frequent itemsets will be reduced.

Disadvantage:

- The second phase requires a full scan of frequent 2-itemsets that would be used when building the graphical structure.
- (5). MFIPA Algorithm

Mining Frequent Itemsets Based On Projection Array (MFIPA) algorithm [8] of mining frequent itemsets based on projection array. The data structure PArray is generated by using data horizontally and vertically firstly, it stores all the frequent 1-itemsets and those items that co-occurrence with signal frequent item. And the usage of PArray avoids the generation of large numbers of candidate itemsets and redundant detection.

Theorem 1. Let PArray[j].item be a frequent itemset with sup(PArray[j].item) = minsup, then there exists no item PArray[j].item < i and i $\notin PArray[j]$.projection, such that PArray[j].item U is a frequent itemset.

Theorem 2. Let item be a single item, the projection of it is $projection(item) = a_1, a_2, \dots, a_m$, if item is combined with the 2^m -1 nonempty subsets of a_1, a_2, \dots, a_m , the supports of the resulting itemsets are all sup(item).

According to Theorem 1 and 2, find some frequent itemsets which have the same support as single frequent item by connecting the single frequent item with every nonempty subsets of its projection, then the frequent itemsets are extended by using depth-first search strategy when the support of single frequent item is greater than minsup. The support of resulting itemsets after extending can be figured out by intersection operation of *tidset(items)*, if the support is greater than minsup, the resulting itemsets will be reserved, and on the contrary, the resulting itemsets will be pruned.

III. Comparision Of Different Algoritms For Fpm

- All the algorithms produce frequent itemsets on the basis of minimum support.
- Apriori works well in sparse data sets since most of the candidates that Apriori generates turn out to be frequent patterns.
- For FP-Tree performs better than all discussed above algorithms because of no generation of candidate sets but the pointes needed to store in memory require large memory space.
- DHP works well at early stages and performance deteriorates in later stages and also results in I/O overhead. In DHP algorithm colliding problem is encounter. At the same bucket two different entry is encountered. It takes a lot of time to scan each transaction of the database for many times.
- Execution time of the first pass of DHP is slightly larger than Apriori due to extra overhead required for generating H₂DHP incurs significantly smaller execution times than Apriori in later passes.
- Apriori Scans full database for every pass whereas DHP scan full database for early two passes and then scan reduced database Dk thereafter.
- For FIG algorithm is a quick mining process that does not use candidates but requires a full scan of frequent 2-itemsets that would be used when building the graphical structure in second phase.

• For MFIPA algorithm works well for dense datasets but performance decreases for sparse datasets[7].

Therefore these algorithms are not sufficient for mining the frequent itemsets for large transactional database.

IV. Proposed New Procedure

The k-dimensional vector is employed to stand for k-itemset, as for the column content which need to be put into a column address, if the column address to contain it had been existed, this column is called old column,

we put the content into corresponding old column address, at the same time, the column count increased by 1; if the column address to contain it had not been existed, a new column address is added to contain it, at the same time, the column count of the new column address is 1.

New Procedure:

Denote the set of existed column address as A, firstly let $A=\Phi$.

Scan the transaction database, take out the transaction one by one, when got the 2-itemset $\{i_x, i_y\}$ from the use 2-dimensional vector (x, y) as the column address of $\{i_x, i_y\}$, let transaction that is being dealt with, n_{xy} be the count of $\{i_x, i_y\}$, check A, if $(x, y) \notin A$, then let $(x, y) \in A$ and $n_{xy} \notin 1$; then let $n_{xy} = n_{xy} + 1$. Repeat steps 2 until all transaction have been dealt with, then we obtain a table. According to prespecified minimum support, frequent itemsets are generated.

Example:

Take the data in Table 3 as an example, we will show how new procedure works.

Table 3: Transactional Data					
ID	List of item_Ids				
1	i_1, i_2, i_5				
2	i_2, i_4				
3	i_2, i_3				
4	i_1, i_2, i_4				
5	i_1, i_3				
6	i_2, i_3				
7	i_1, i_3				
8	i_1, i_2, i_3, i_5				
9	i_1, i_2, i_3				

Set $A = \Phi$.

Scan the transaction database, take out the transaction one by one. As for transaction 1, generate the following 2-itemsets: $\{i_1, i_2\}, \{i_1, i_5\}$ and $\{i_2, i_5\}$. As for $\{i_1, i_2\}$, since $(1, 2) \notin A$, then let $A = \{(1, 2)\}$ and $n_{12} = 1$. As for $\{i_1, i_5\}$, since $(1,5) \notin A = \{(1,2)\}$, then let $A = \{(1,2), (1,5)\}$ and $n_{15} = 1$. As for $\{i_2, i_5\}$, since $(2,5) \notin A = \{(1, 2), (1, 5)\}$ 2),(1,5), then let $n_{25} = 1$ and $A = \{(1, 2), (1, 5), (2, 5)\}.$

As for transaction 2, generate the 2-itemset $\{i_2, i_4\}$, as for $\{i_2, i_4\}$, since $(2, 4) \notin A = \{(1, 2), (1, 5), (2, 5)\}$, then $A = \{(1,2), (1,5), (2,5), (2,4)\}$ and let $n_{24} = 1$.

As for transaction 3, generate the 2-itemset $\{i_2, i_3\}$, as for $\{i_2, i_3\}$, since $(2,3) \notin A = \{(1,2), (1,5), (2,5), (2,4)\}$, then A = {(1, 2), (1,5), (2,5), (2,4), (2,3)} and $n_{23} = 1$. let

As for transaction 4, generate the following 2-itemsets: $\{i_1, i_2\}, \{i_1, i_4\}$ and $\{i_2, i_4\}$. As for $\{i_1, i_2\}$, since (1, 2) \in A, then let $n_{12} = 1 + 1 = 2$. As for $\{i_1, i_4\}$, since $(1, 4) \notin A = \{(1, 2), (1, 5), (2, 5), (2, 4), (2, 3)\}$, then let $A = \{(1, 2), (1, 5), (2, 5), (2, 4), (2, 3)\}$, then let $A = \{(1, 2), (1, 5), (2, 5), (2, 4), (2, 3)\}$. $\{(1,2),(1,5),(2,5),(2,4),(2,3),(1,4)\}$ and $n_{14} = 1$. As for $\{i_2, i_4\}$, since $(2,4) \in A = \{(1,2), (1,5), (2,5), (2,4), (2,5), (2,4), (2,5), (2,5), (2,4), (2,5),$ (2,3),(1,4), then let $n_{24} = 1 + 1 = 2$.

As for the other transactions and 2-itemsets, we can deal with them similarly. Finally, we obtain a table shown in Table 4.

Address	(1, 2)	(1, 5)	(2, 5)	(2, 4)	(2, 3)	(1, 4)	(1, 3)	(3, 5)
Count	4	2	2	2	4	1	4	1
Content	{i1, i2}	{i1, i5}	{i2, i5}	{i2, i4}	{i2, i3}	{i1, i4}	{i1, i3}	{i3, i5
	{i1, i2}	{i1, i5}	{i2, i5}	{i2, i4}	{i2, i3}		{i1, i3}	
	{i1, i2}				{i2, i3}		{i1, i3}	
	$\{i1, i2\}$				$\{i2, i3\}$		{i1, i3}	

Table 4: Table using by New Procedure

V. Conclusion

New procedure has some advantages. Its rationale is simple and clear, it can be carried out without any Hash function and will not encounter Hash colliding problem, it absorbed the merits of 2-dimensional Hash algorithm which employ 2-dimensional vector to represent column address. New procedure generate the appropriate column address just for the column content, the number of address is the most but necessary and the least but sufficient number needed to avoid Hash colliding, that is it is the best value.

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